

PrivateDrop: Practical Privacy-Preserving Authentication for Apple AirDrop



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ENCRYPTO
CRYPTOGRAPHY AND
PRIVACY ENGINEERING



SEMG
SECURE MOBILE NETWORKING

privatedrop.github.io

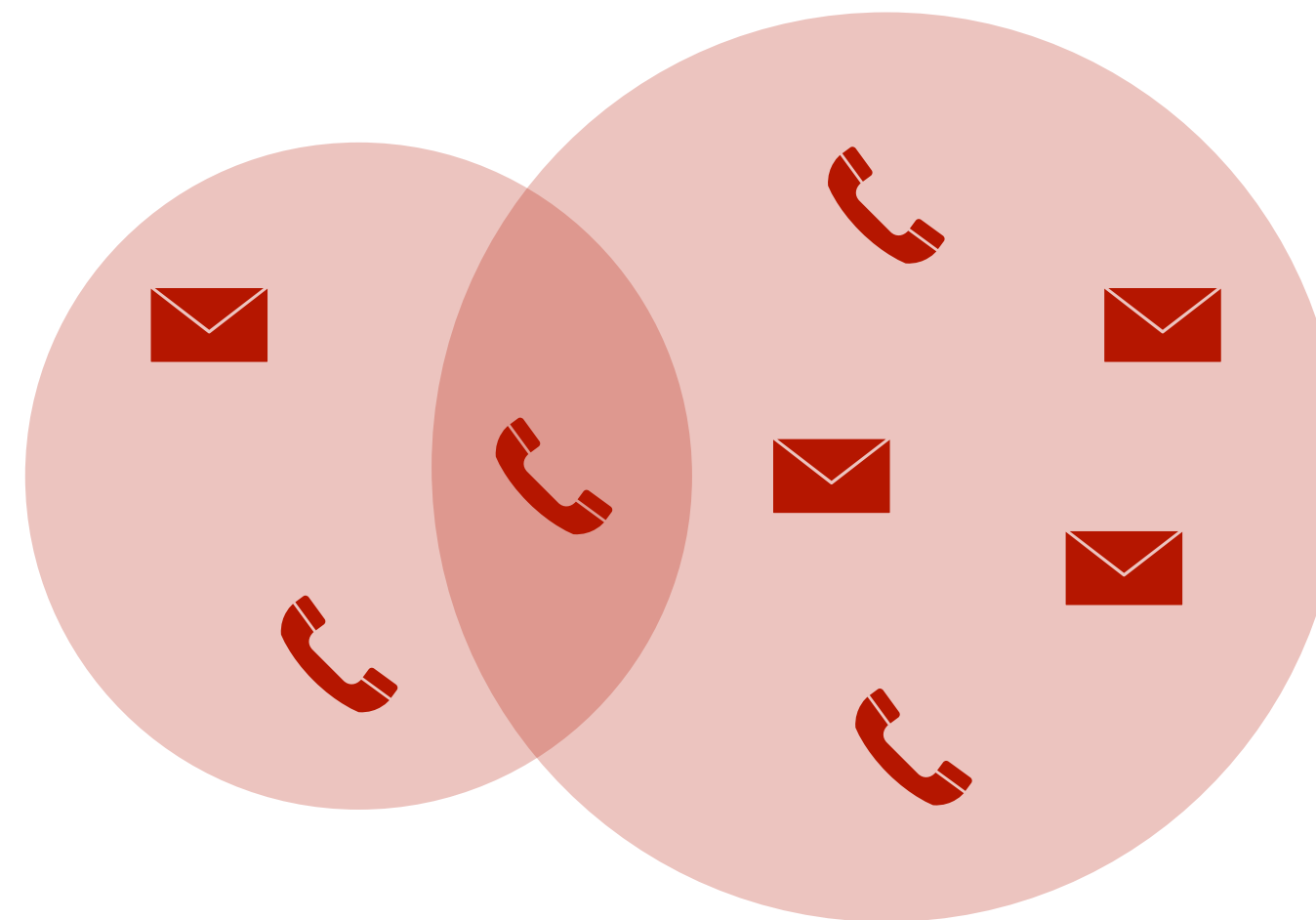


Roadmap to PrivateDrop



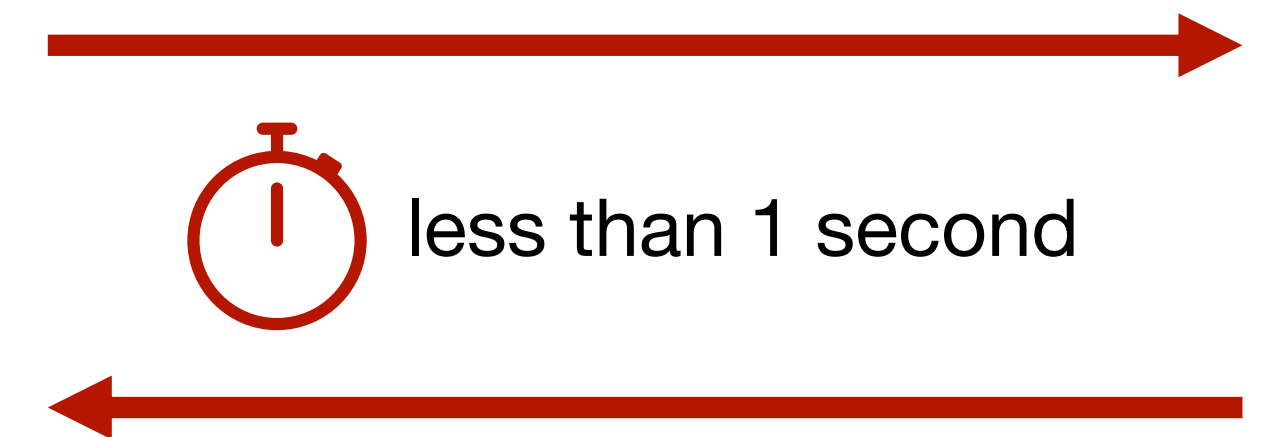
Discover

Contact Identifier Leakage
by Apple AirDrop



Design

Privacy-Preserving Authentication
via Private Set Intersection



Demonstrate

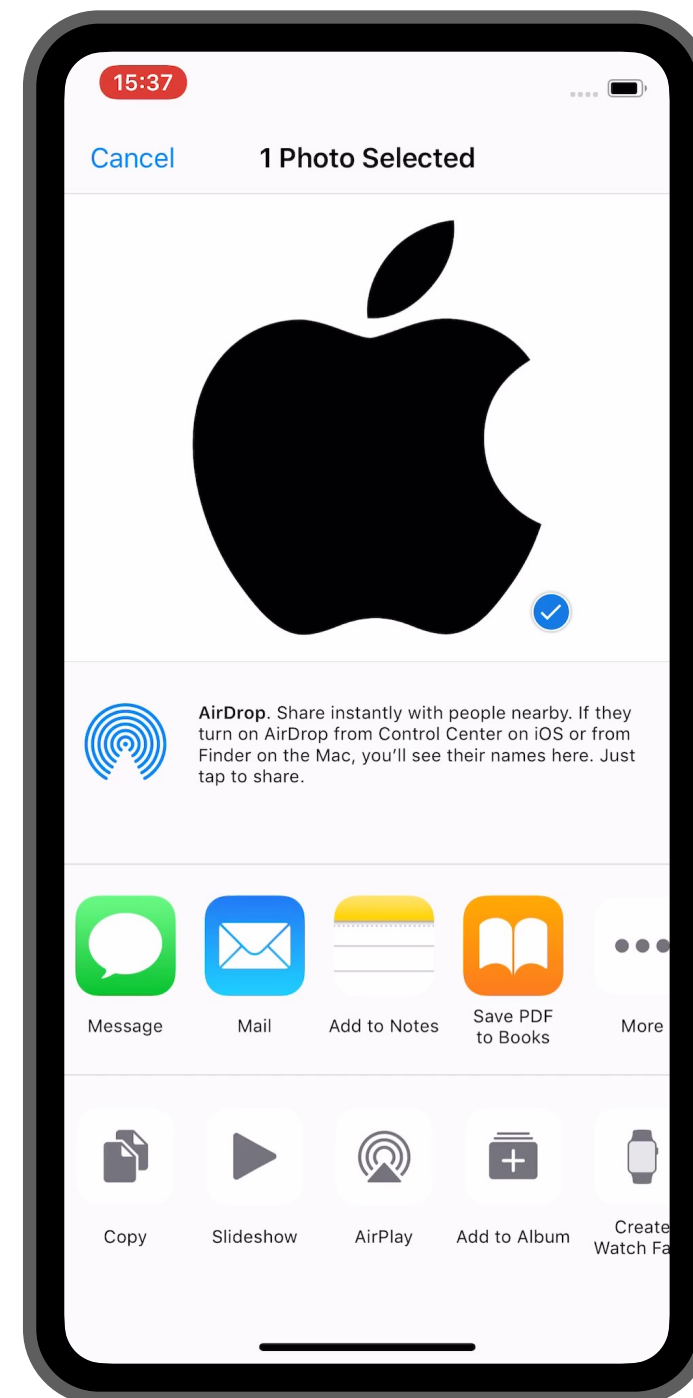
Native Prototype
with Excellent User Experience

AirDrop Authentication

[SNMHKNH19]

Find out whether we are mutual contacts

via Wi-Fi/AWDL [SKH18]

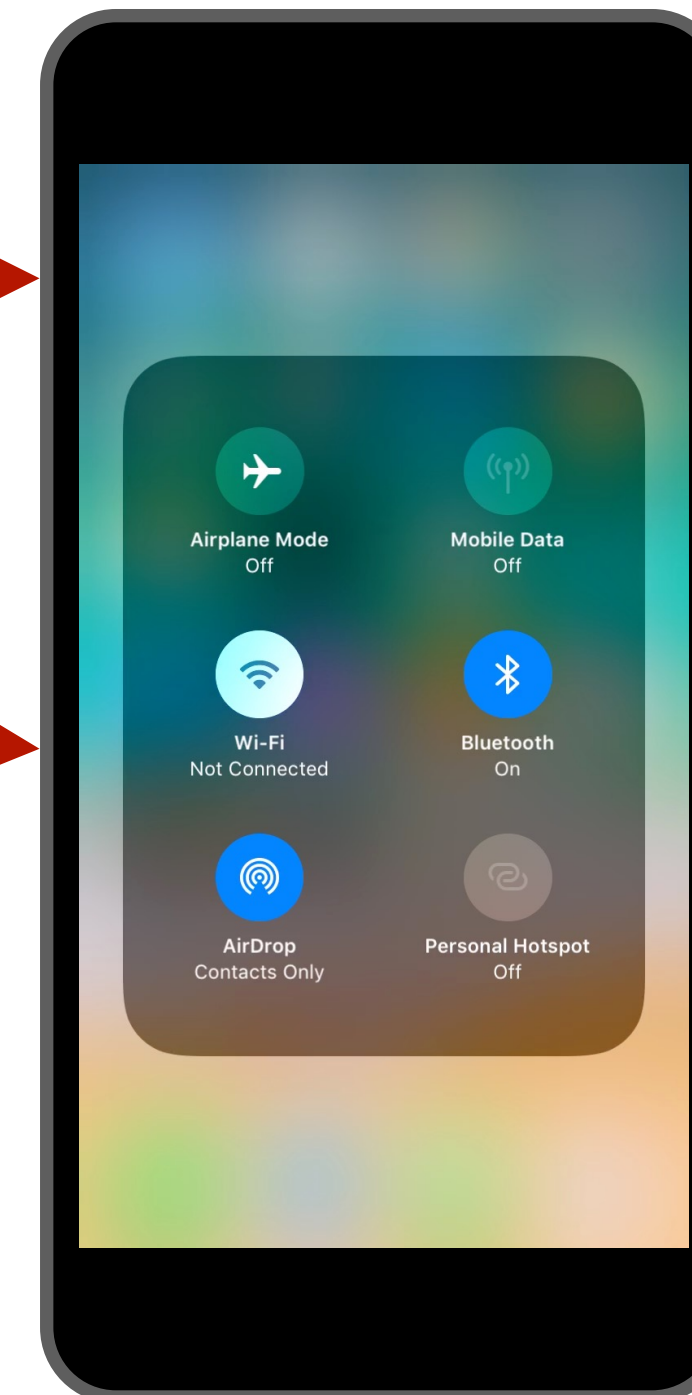


Sender

TLS connection with client and server certificates

HTTP POST /Discover with sender's validation record*

I want to find other people, so I tell them who I am

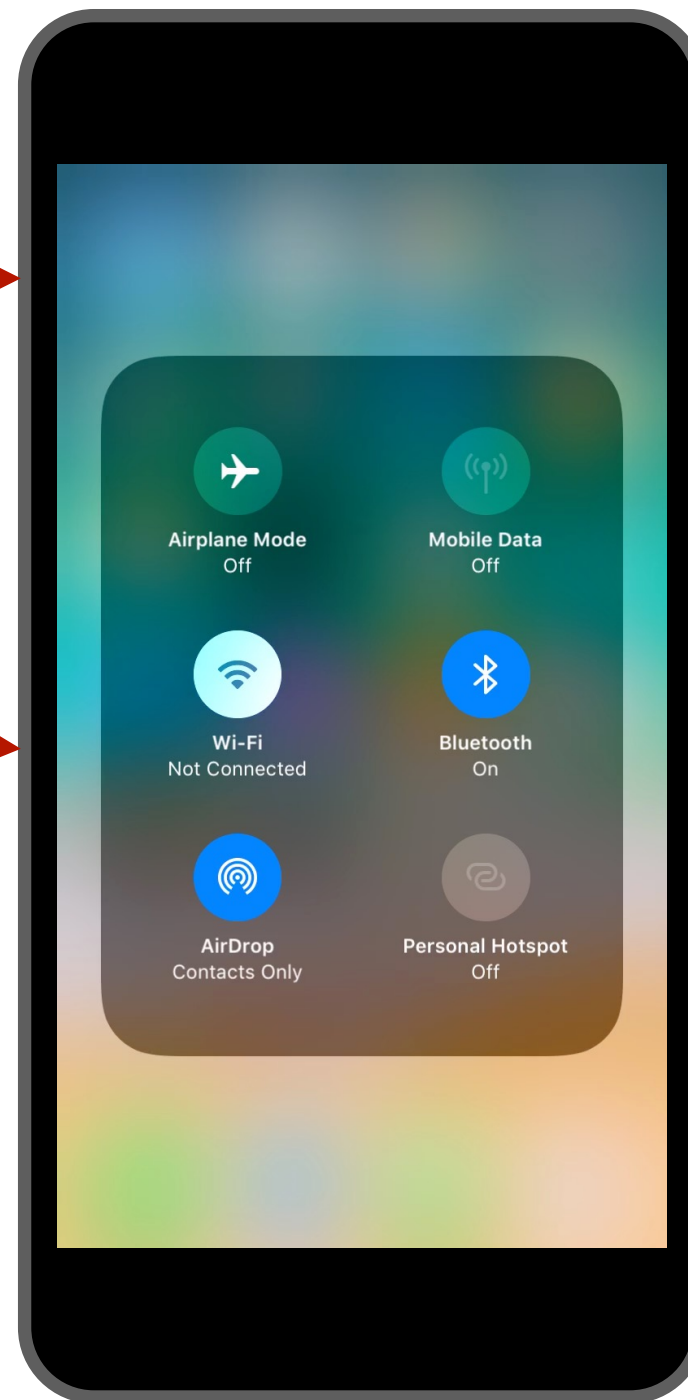
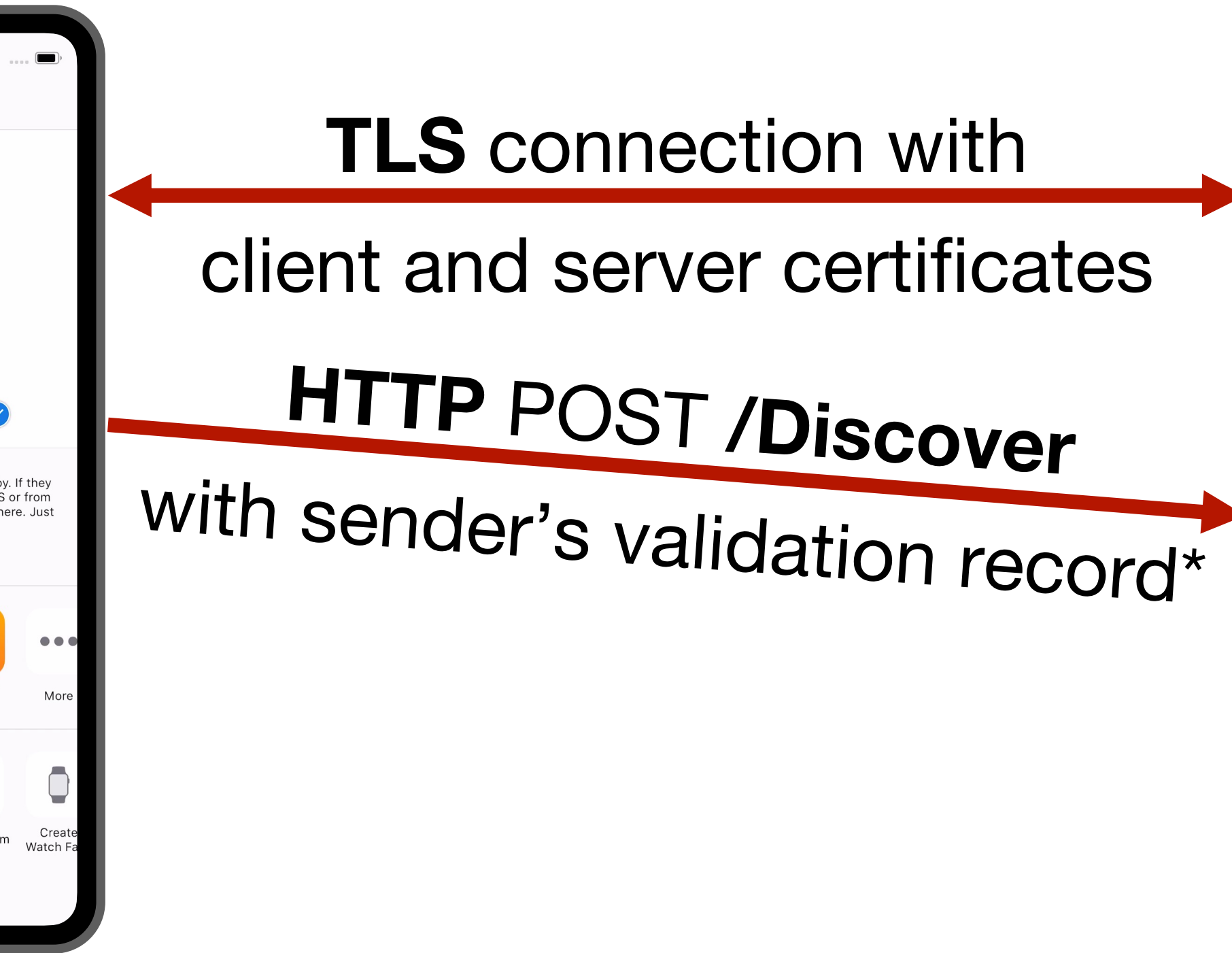


Receiver

* Apple-signed cert including
 $H_i = SHA256(+49\ 123\ \dots)$
 $H_j = SHA256(\dots\ @icloud.com)$

AirDrop Authentication

[SNMHKNH19]



Receiver

$\exists H_i \in VR: H_i \in \textit{address book}$
(+ check validation record
+ check TLS certificate)

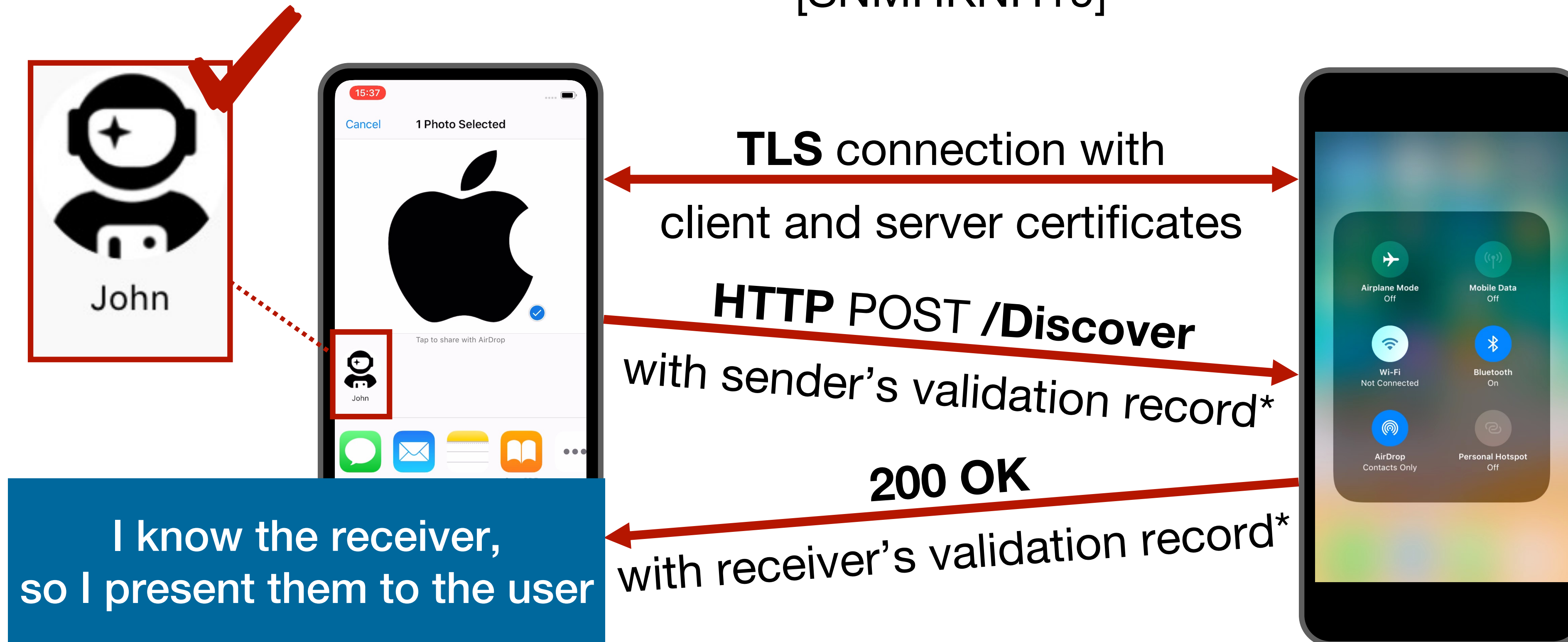


I know the sender,
so I tell them who I am

* Apple-signed cert including
 $H_i = \textit{SHA256}(+49\ 123\ \dots)$
 $H_j = \textit{SHA256}(\dots\ @\textit{icloud.com})$

AirDrop Authentication

[SNMHKNH19]



Sender

Receiver

* Apple-signed cert including
 $H_i = \text{SHA256}(+49\ 123\ \dots)$
 $H_j = \text{SHA256}(\dots\ @\text{icloud.com})$

AirDrop Authentication: What can go wrong here?

Sender

I want to find other people,
so I tell them who I am



“Hashing vs. Hiding”
(Sender Leakage)

$$H_{S,i} = \text{SHA256}(+49\ 123\dots)$$

Receiver



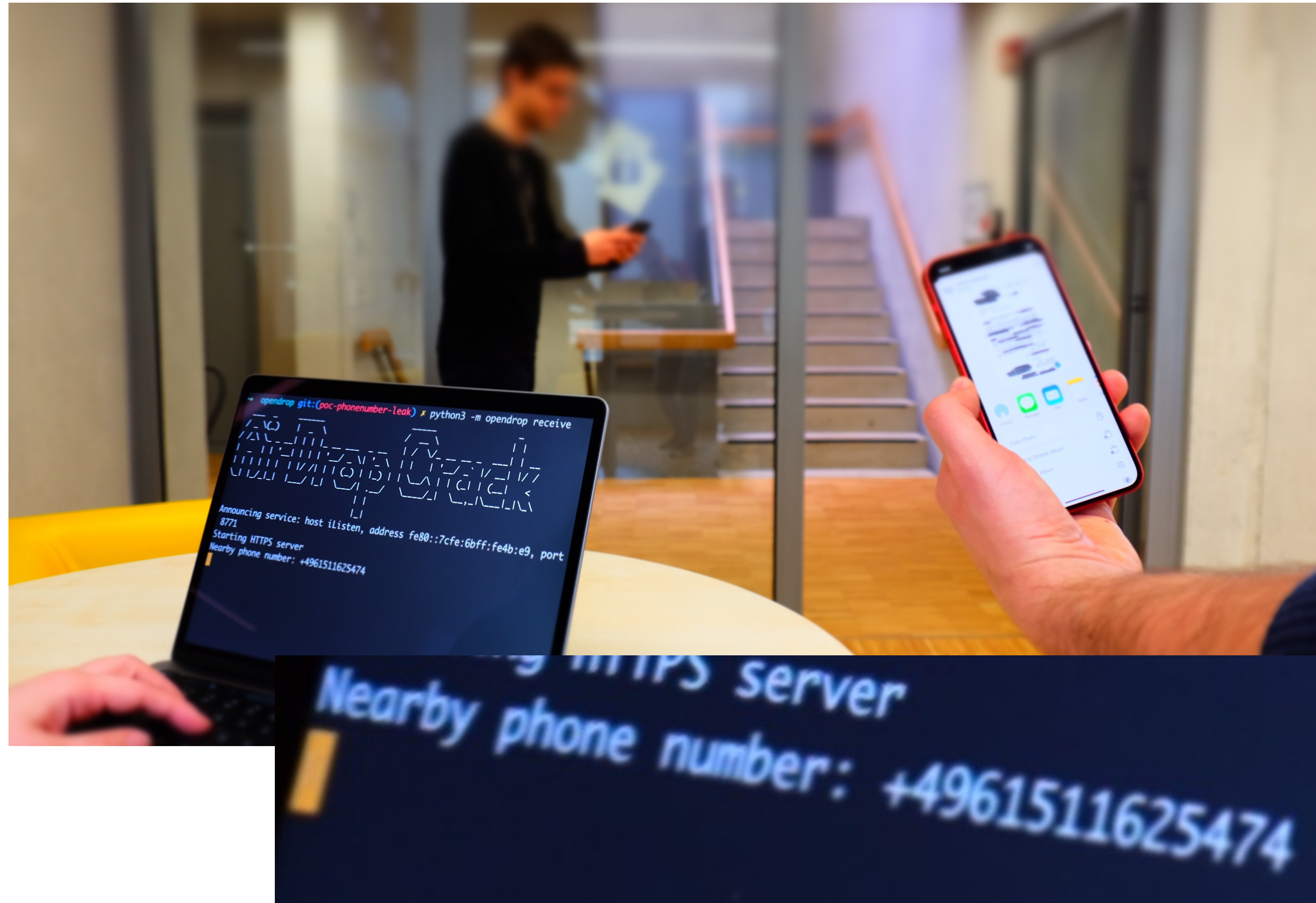
“Celebrity Issue”
(Receiver Leakage)

I know the sender,
so I tell them who I am

$$H_{R,i} = \text{SHA256}(+1\ 234\dots)$$

I know the receiver,
so I present him to the user

Exploiting the Vulnerabilities in Practice



Requirements

- Physical proximity to target
- Wi-Fi-capable device

Proof-of-concept

- “AirCollect”
<https://github.com/seemoo-lab/opendrop/blob/poc-phonenum-leak/README.PoC.md>
- Makes use of optimized rainbow tables [HWSDS21]

Impact

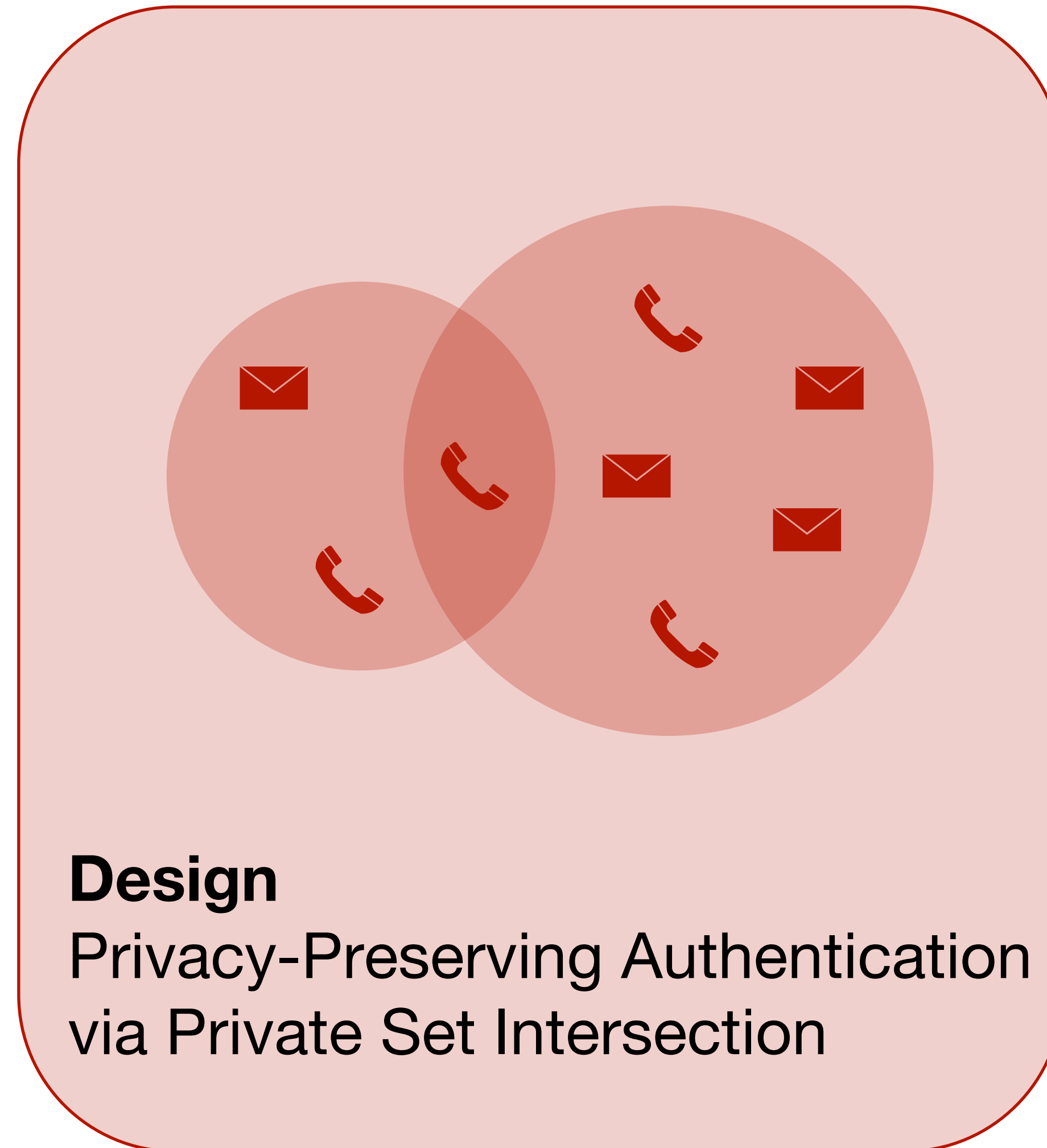
- Recover phone numbers of AirDrop users in real-time

Roadmap to PrivateDrop



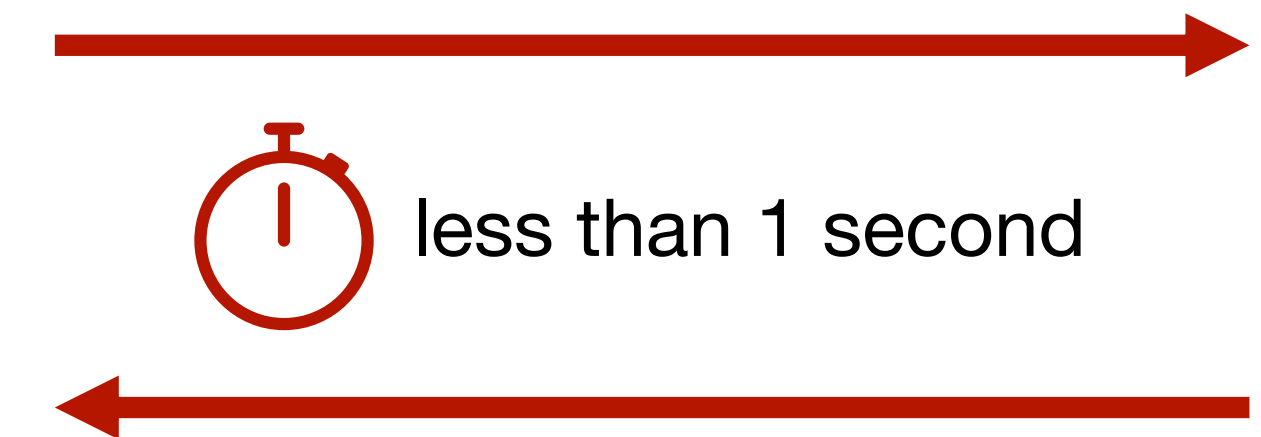
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Design

Privacy-Preserving Authentication
via Private Set Intersection



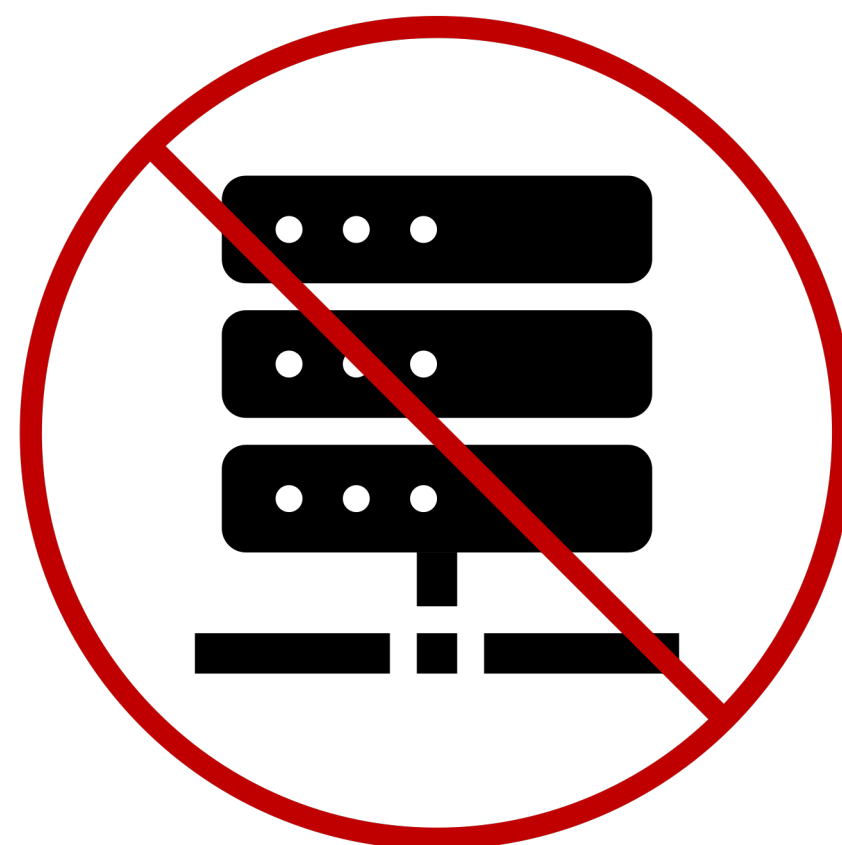
Demonstrate

Native Prototype
with Excellent User Experience

PrivateDrop Requirements

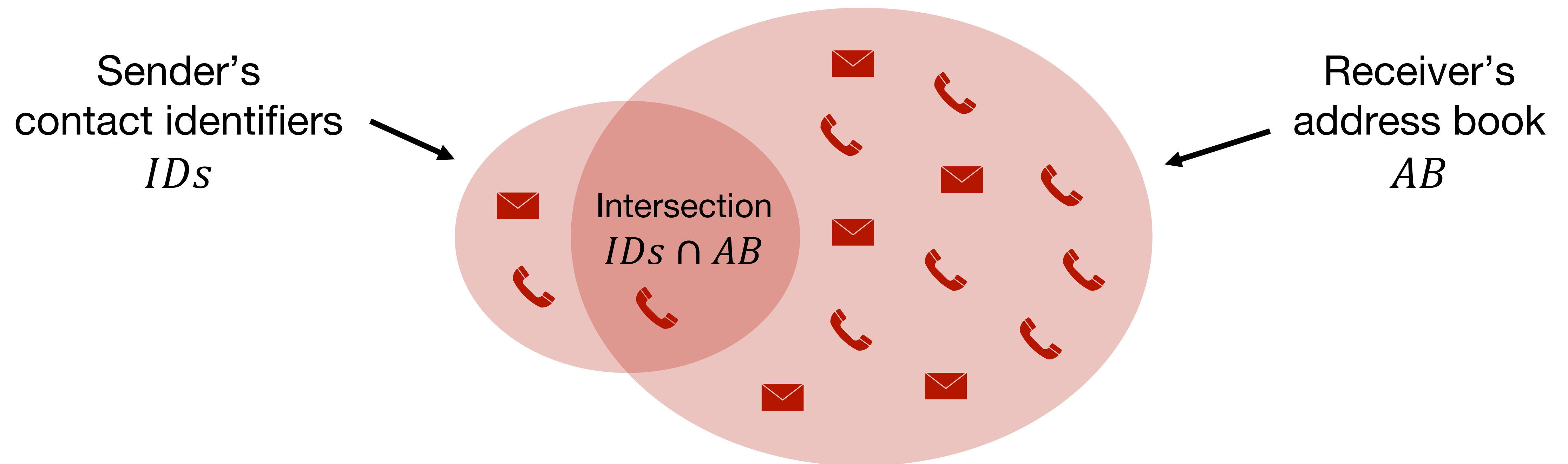
Privacy requirements:

1. Disclose contact identifiers only *if both parties are mutual contacts*.
2. Only disclose those *contact identifiers that the other party already knows*.



Apply private set intersection (PSI) to achieve private mutual authentication

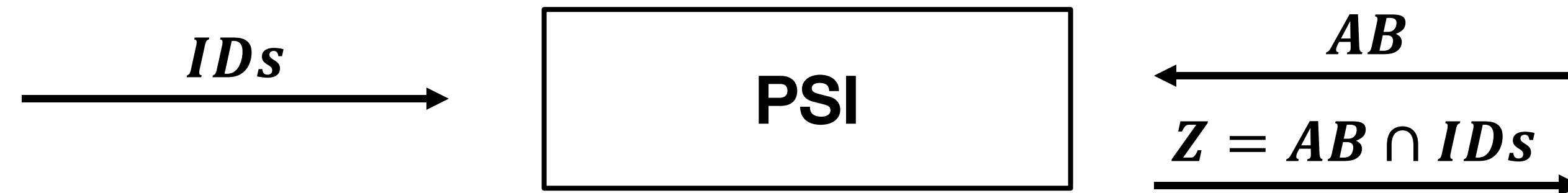
Private Set Intersection (PSI)



AirDrop: Semantics

AirDrop Sender S

AirDrop Receiver R



“I know S ”



“I know R ”

Problems:

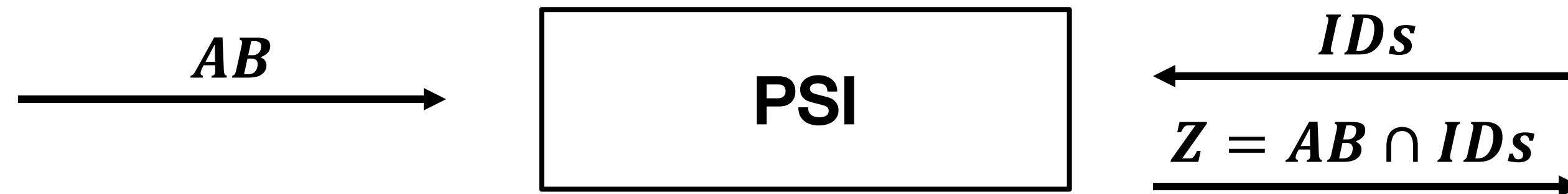
- Malicious receivers
- Online complexity depends on AB (large)

AB: address book
IDs: contact identifiers

PrivateDrop: Semantics

AirDrop Sender S

AirDrop Receiver R



“S knows me”



“R knows me”

Next: S and R can disclose their known identities, i.e., $IDs \cap AB$

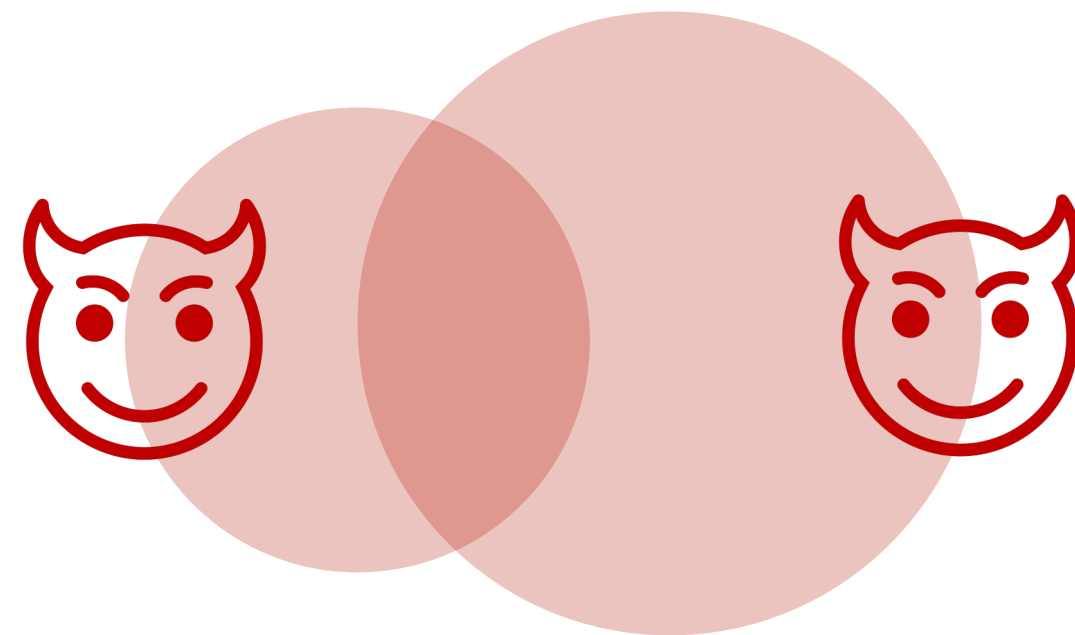
Changed Semantics:

- Receivers in check
- Online complexity depends on IDs (small)

AB: address book
IDs: contact identifiers

PrivateDrop Design and Implementation

Maliciously Secure
PSI Protocol



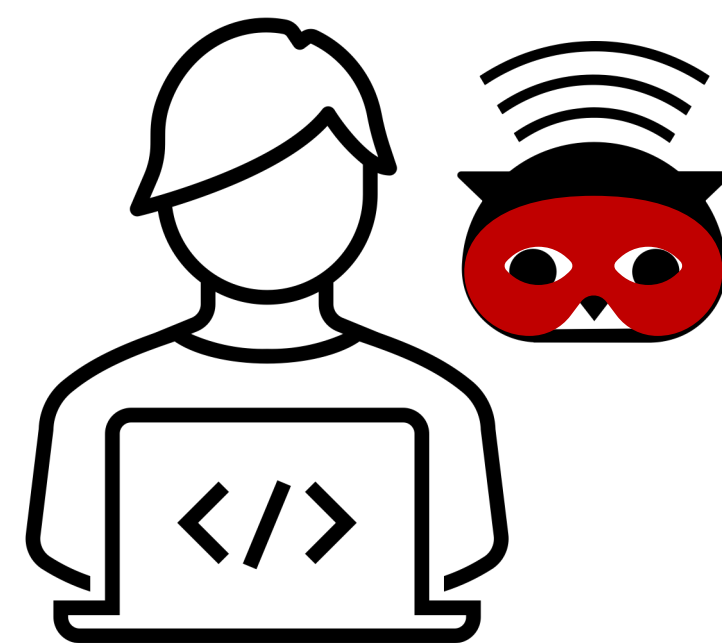
Protection against
Malicious Inputs



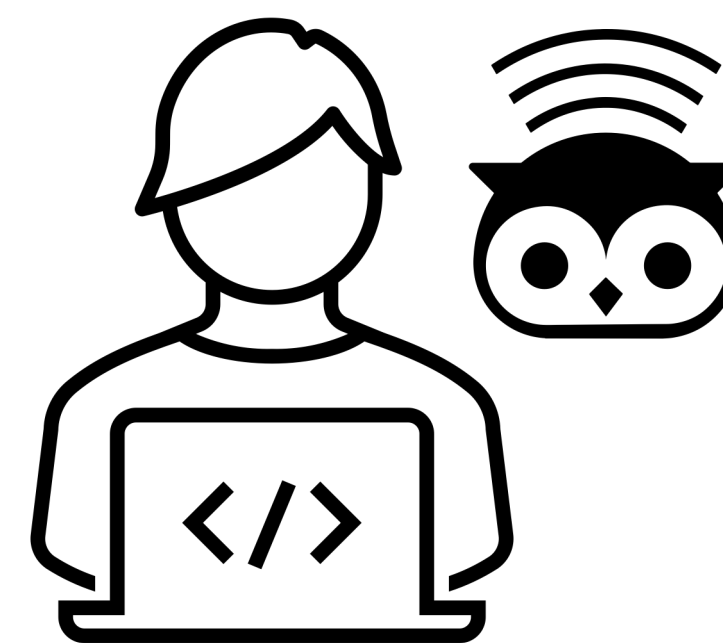
Integration of PSI
into AirDrop



Backwards
Compatibility

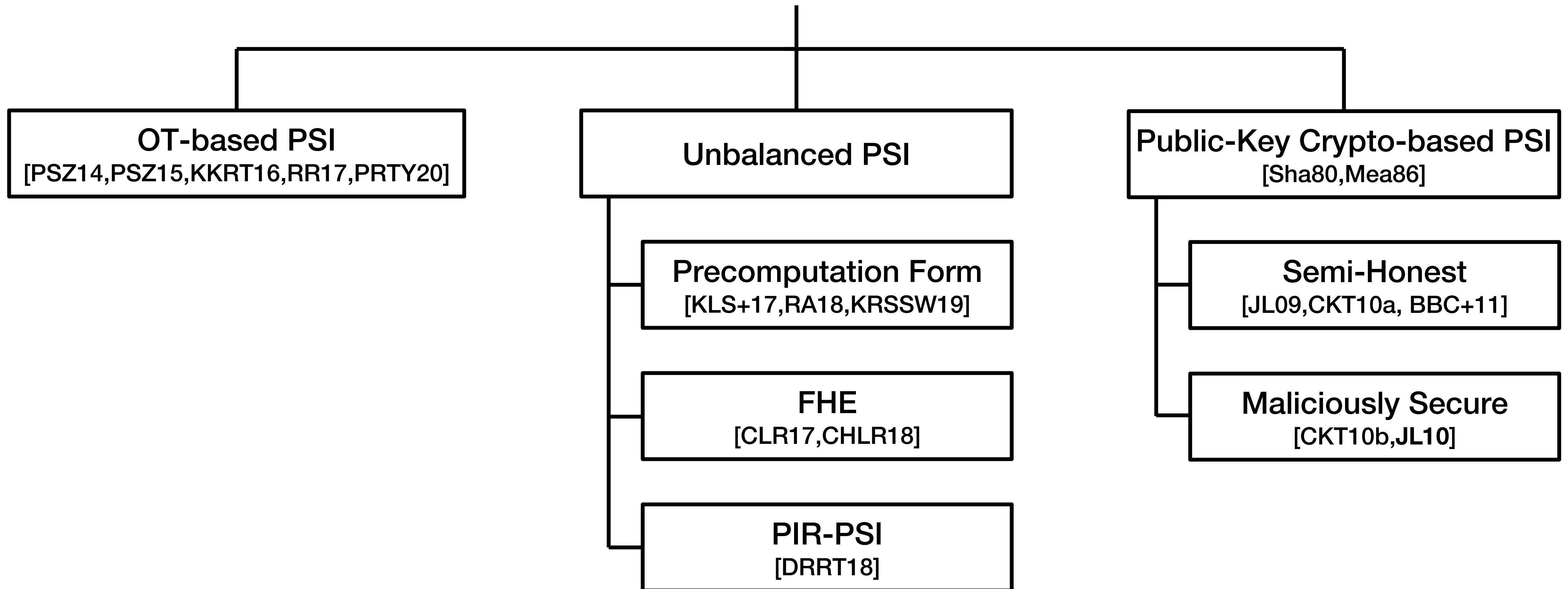


PrivateDrop
Implementation



AirDrop
Implementation

Choice of PSI Protocol



Optimized PSI Protocol of [JL10]


AirDrop Sender S


AirDrop Receiver R



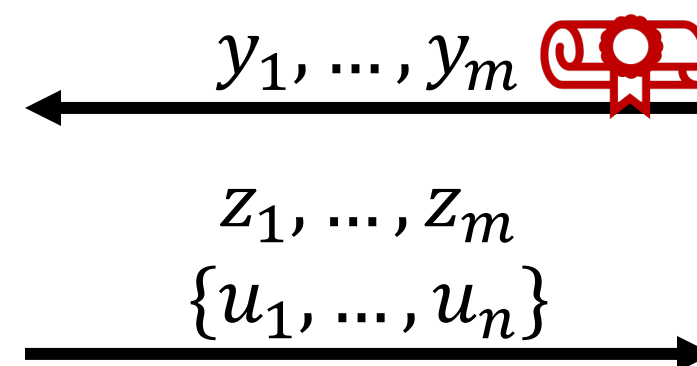
$k \xleftarrow{\$} \mathbb{Z}_q$
 For $j = 1$ to n :
 $u = H(H(c_j), H(c_j)^k)$

Precomputation

For $j = 1$ to m :
 $\alpha_i \xleftarrow{\$} \mathbb{Z}_q$
 $h_i = H(ID_i)$
 $y_i = (h_i)^{\alpha_i}$ Obtain  for y_i

For $i = 1$ to m :
 Verify 
 $z_i = y_i^k$

Online



For $i = 1$ to m :
 $v_i = H(h_i, (z_i)^{1/\alpha_i})$
 Output $\{ID_i \in IDs \mid \exists j : u_j = v_i\}$

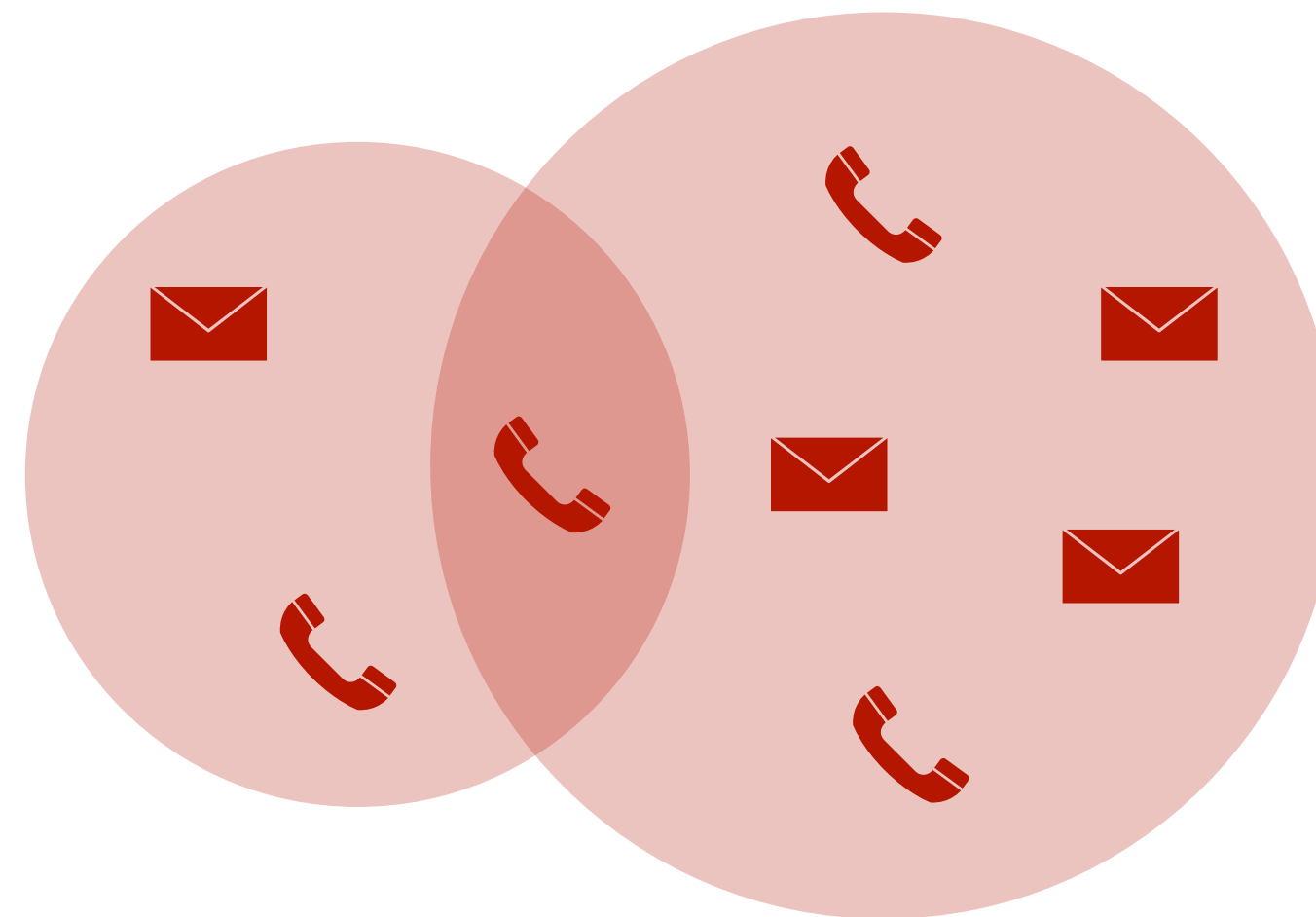
(simplified version, omits ZK proofs for malicious security)

Roadmap to PrivateDrop



Discover

Contact Identifier Leakage
by Apple AirDrop



Design

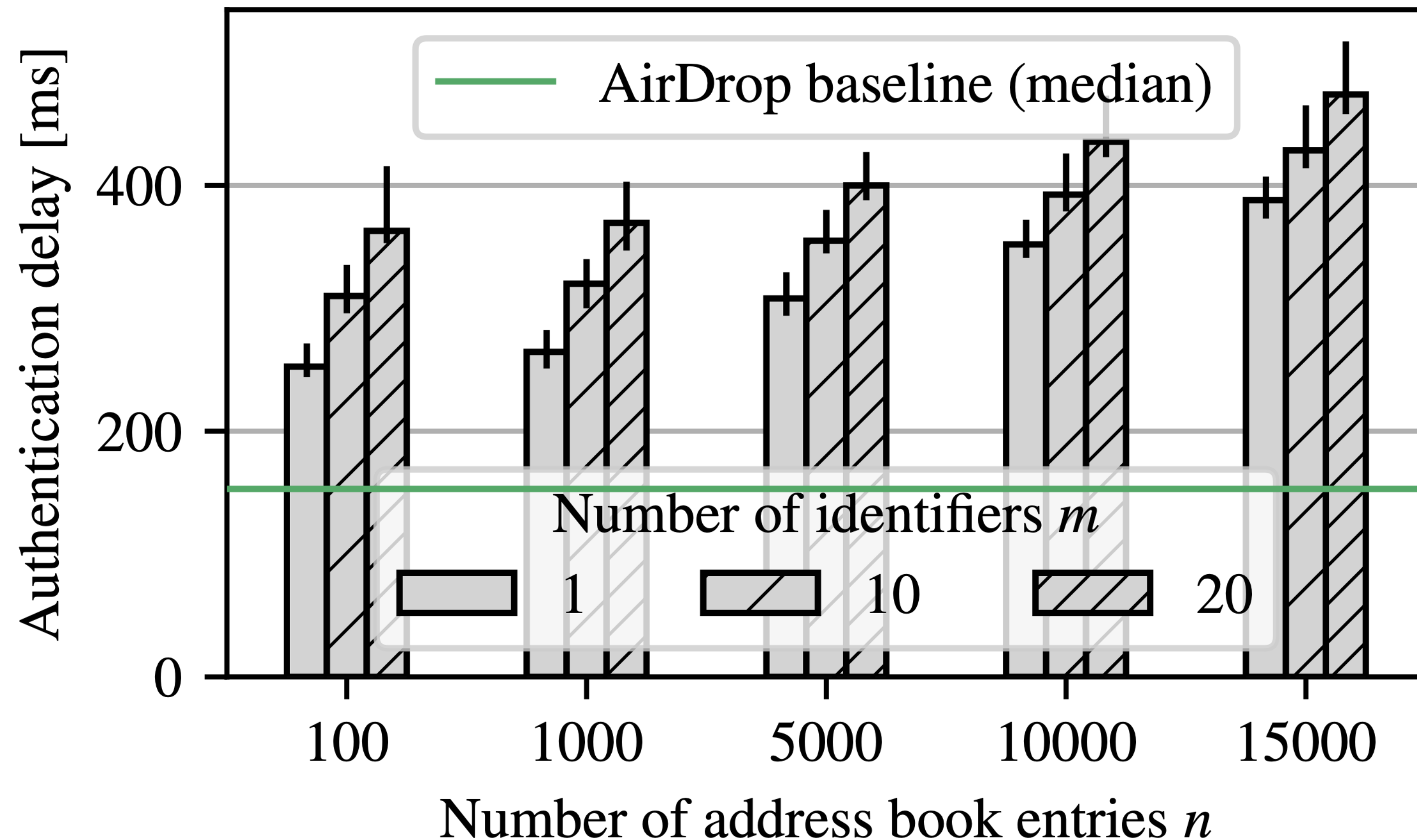
Privacy-Preserving Authentication
via Private Set Intersection

less than 1 second

Demonstrate

Native Prototype
with Excellent User Experience

PrivateDrop Results: Authentication Delay



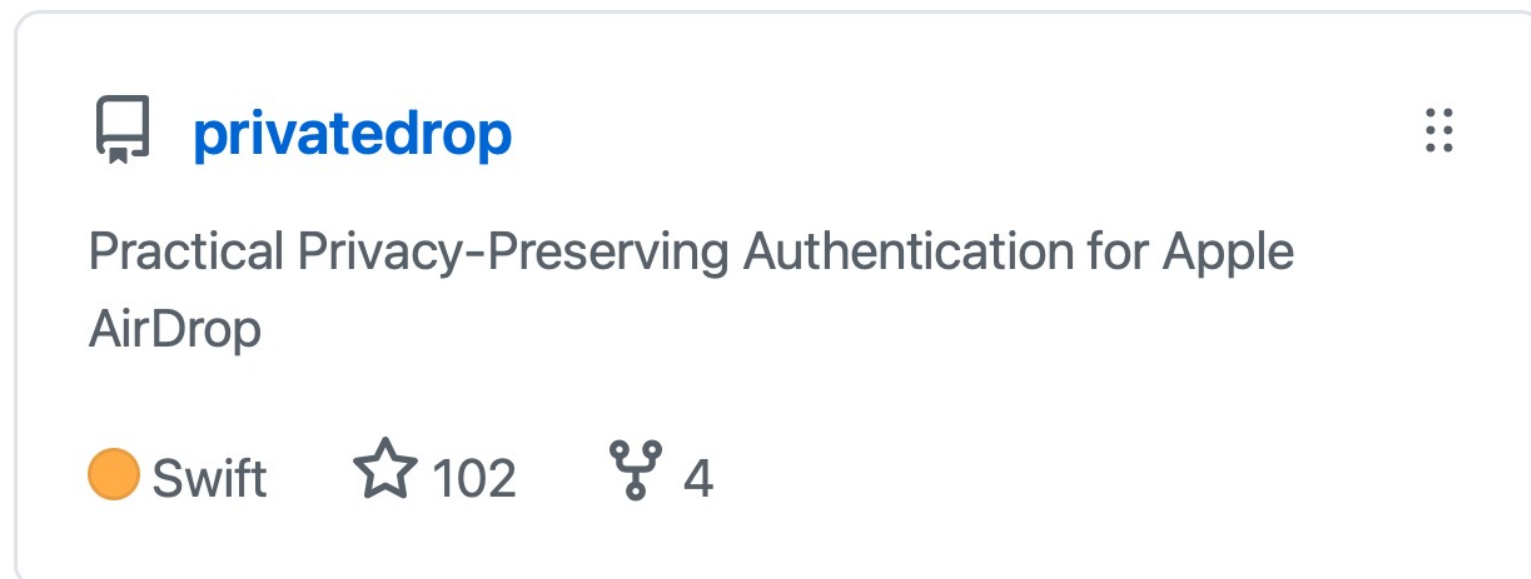
- **Native** implementation for macOS and iOS
- There is **some (expected) overhead**
- But, authentication **delay is well below 1 second** (“immediate response”)

Setup: MacBook and iPhone connected via USB cable (results for Wi-Fi connection with stronger variance in the paper)

PrivateDrop: Privacy-Preserving Mutual Authentication for Apple AirDrop

Open-Source Software

Native implementation for macOS and iOS as open-source software available at privatedrop.github.io



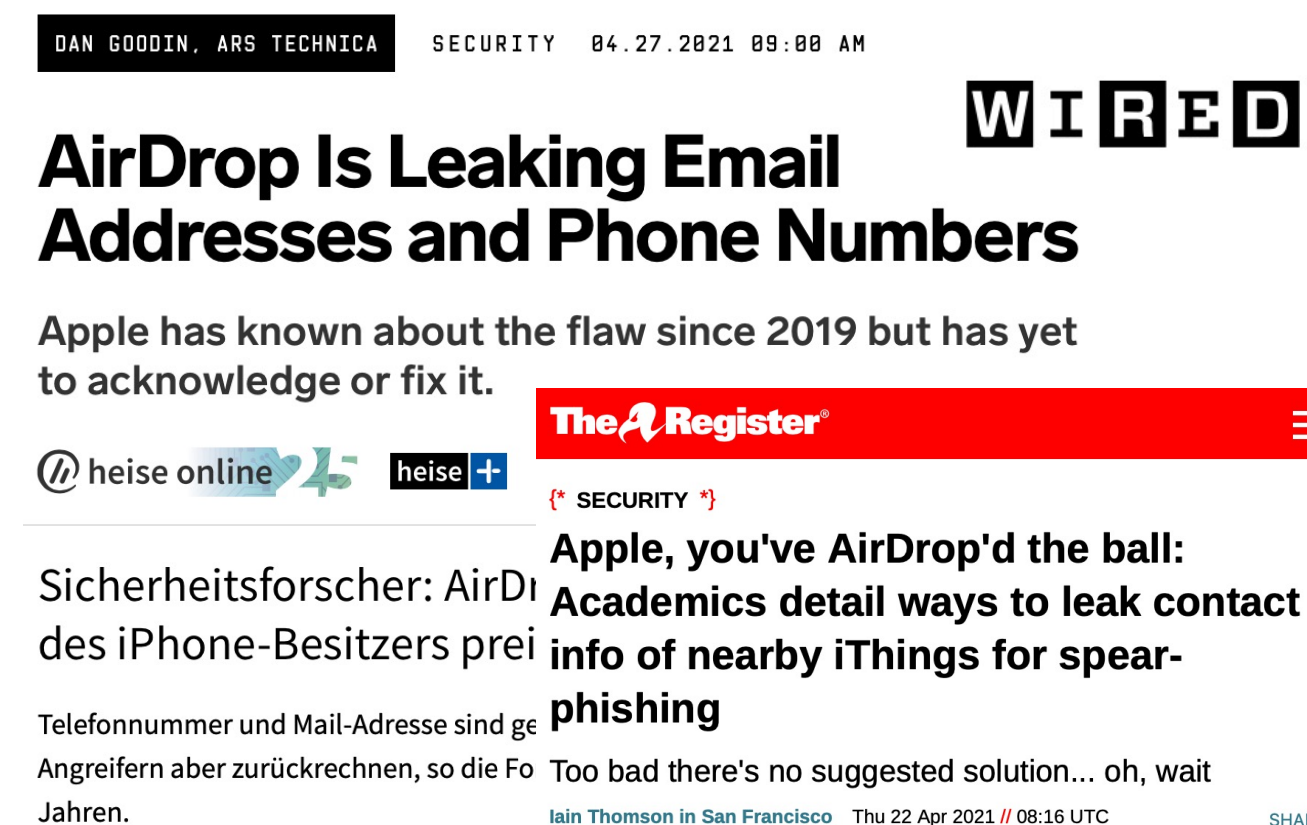
privatedrop

Practical Privacy-Preserving Authentication for Apple AirDrop

Swift 102 4

Press and Media

International and national coverage



DAN GOODIN, ARS TECHNICA SECURITY 04.27.2021 09:00 AM

AirDrop Is Leaking Email Addresses and Phone Numbers

Apple has known about the flaw since 2019 but has yet to acknowledge or fix it.

heise online heise + The Register

SECURITY

Sicherheitsforscher: AirDrop des iPhone-Besitzers preisgibt

Telefonnummer und Mail-Adresse sind gefährlich. Angreifern aber zurückrechnen, so die Forscher in den letzten Jahren.

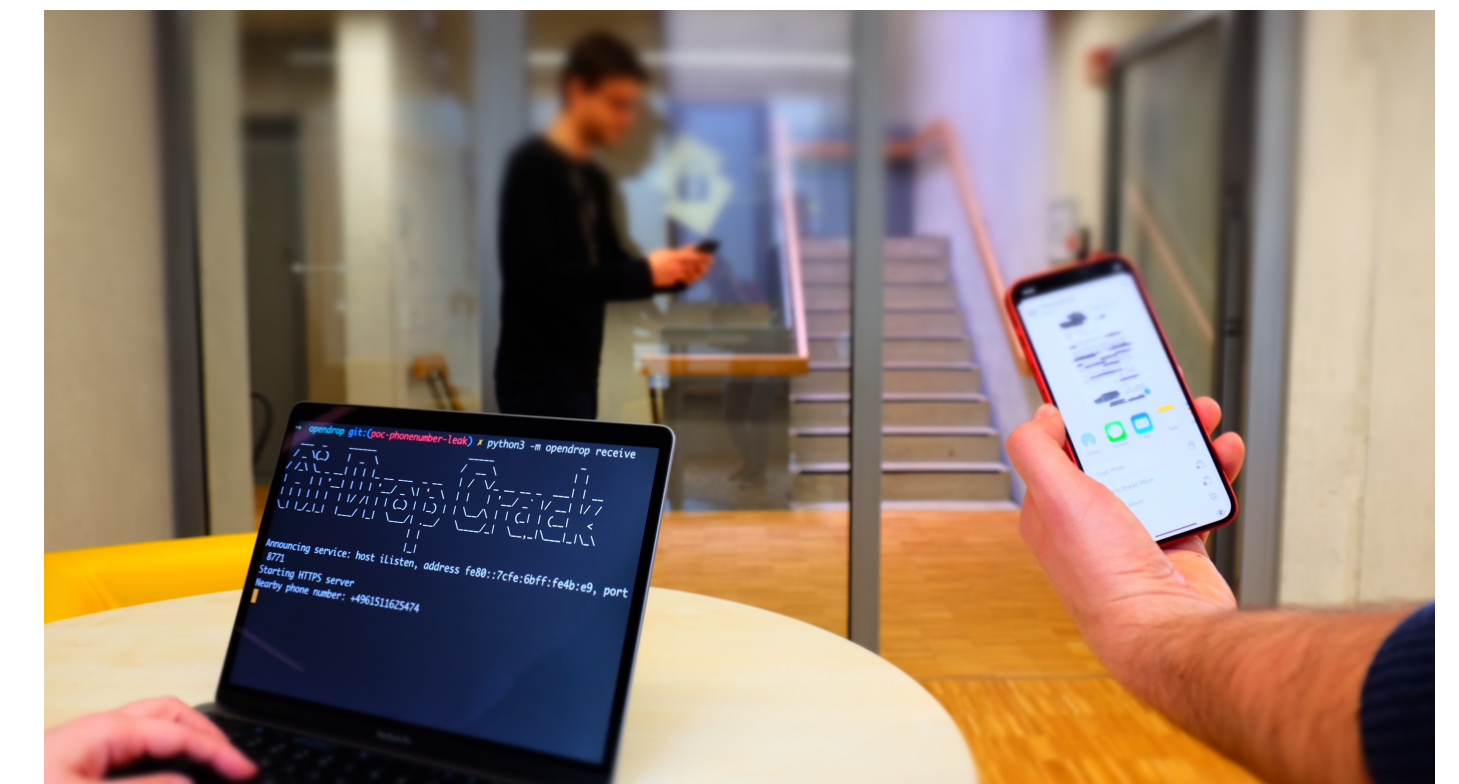
Apple, you've AirDrop'd the ball: Academics detail ways to leak contact info of nearby iThings for spear-phishing

Too bad there's no suggested solution... oh, wait

Iain Thomson in San Francisco Thu 22 Apr 2021 // 08:16 UTC

Responsible Disclosure

Apple users are still vulnerable to the discovered privacy leaks



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Acknowledgements

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