The GNU Name System

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30C3

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Bundesnachrichtendiens

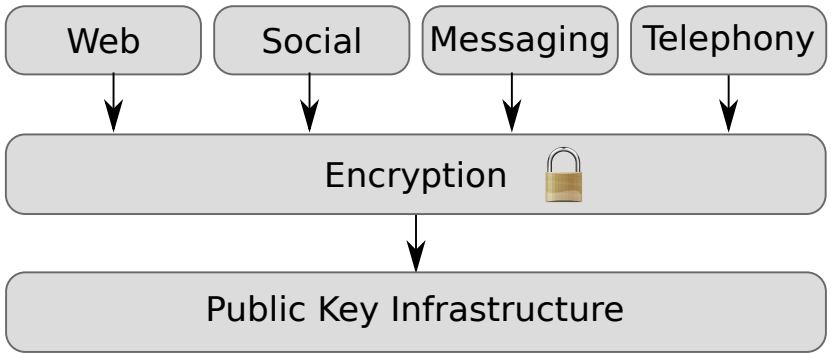
Adversary Model

- Any role
- Multiple Identities
- Computational Power
- Legal Power

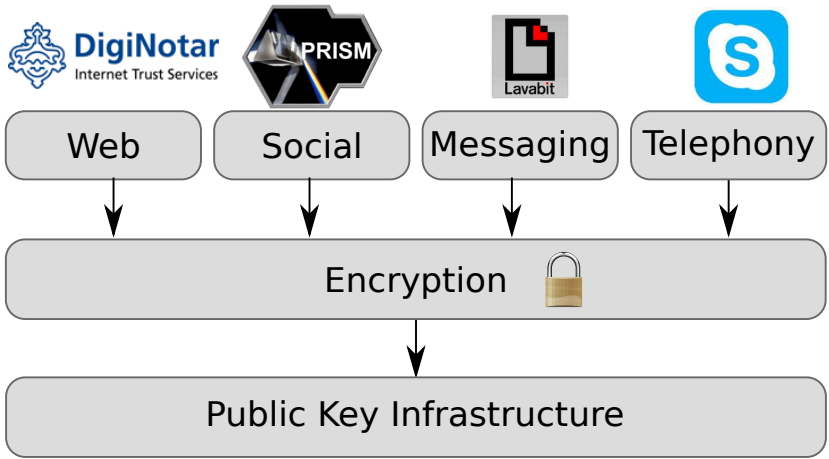
But cannot:

- Break or prevent crypto
- Compromise end-user system
- Prevent network communication

Broken Pillars



Broken Pillars



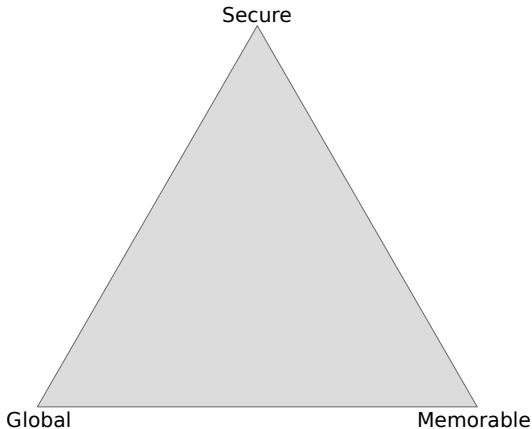
We MUST Decentralize!

Centralized Internet infrastructure is easily controlled:

- Number resources (IANA)
- Domain Name System (Root zone)
- DNSSEC root certificate
- X.509 CAs (HTTPS certificates)
- Major browser vendors (CA root stores!)

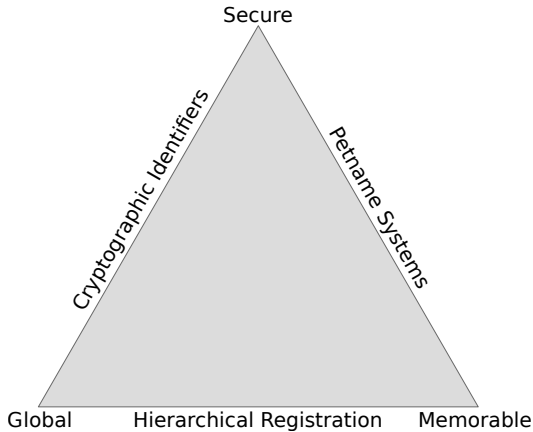
Encryption does not help if PKI is compromised!

Zooko's Triangle



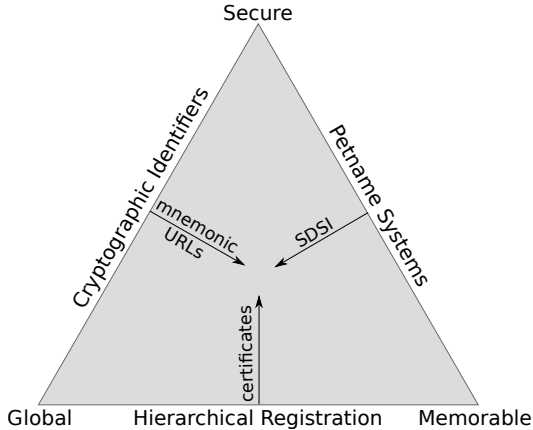
A name system can only fulfill **two**!

Zooko's Triangle



DNS, ".onion" IDs and `/etc/hosts/` are representative designs.

Zooko's Triangle



DNSSEC security is broken by design (adversary model!)

Namecoin

Namecoin

- Memorable:

Namecoin

- Memorable: Check
- Global:

Namecoin

- Memorable: Check
- Global: Check
- Secure:

Namecoin

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- Secure: different adversary model!

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$\Rightarrow$ Availability of names (registration rate) is restricted

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- Memorable: Check
- Global: Check
- Secure: different adversary model!

$\Rightarrow$ Availability of names (registration rate) is restricted

$\Rightarrow$ Adversary must not have 51% compute power

The GNU Name System

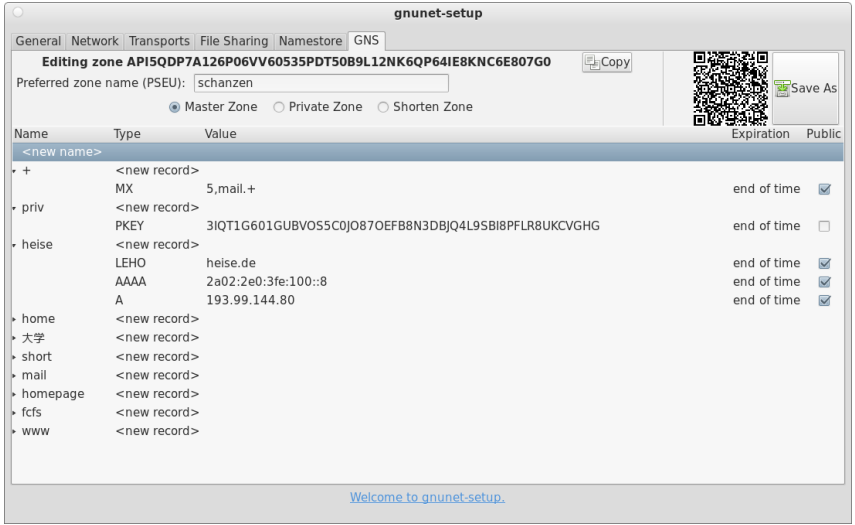
Properties of GNS

- Decentralized name system with secure memorable names
- Delegation used to achieve transitivity
- Also supports globally unique, secure identifiers
- Achieves query and response privacy
- Provides alternative public key infrastructure
- Interoperable with DNS

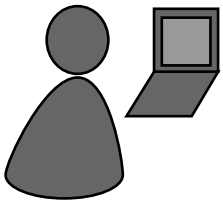
Uses for GNS in GNUnet

- Identify IP services hosted in the P2P network
- Identities in social networking applications

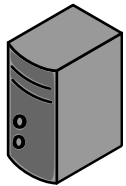
Zone Management: like in DNS



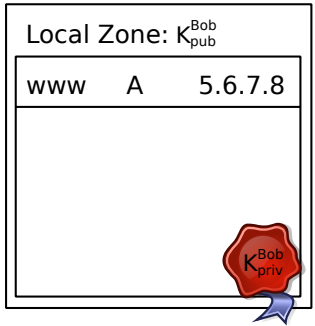
Name resolution in GNS



Bob

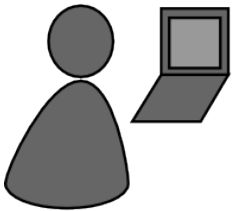


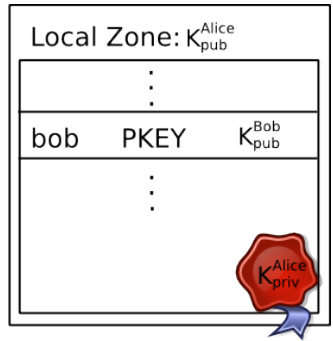
Bob's webserver



Secure introduction

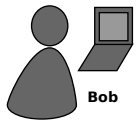
Delegation

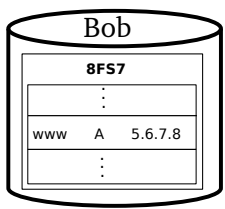


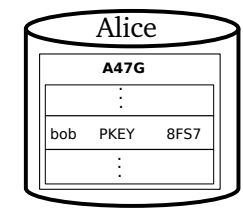


- Alice learns Bob's public key
- Alice creates delegation to zone **bob**
- Alice can reach Bob's webserver via **www.bob.gnu**

Name Resolution







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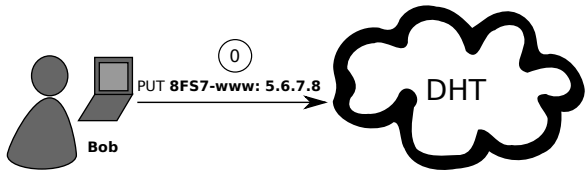

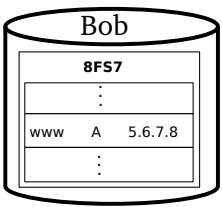

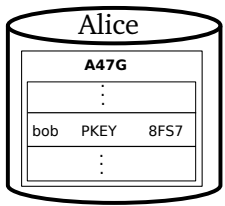

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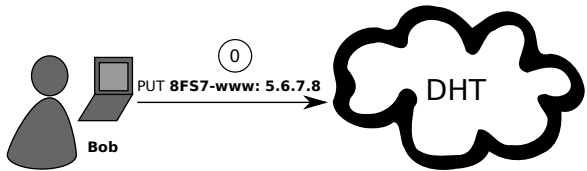

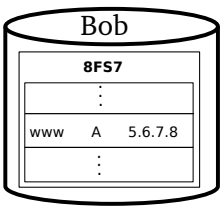

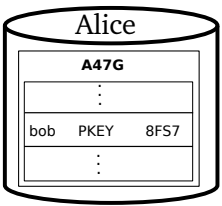

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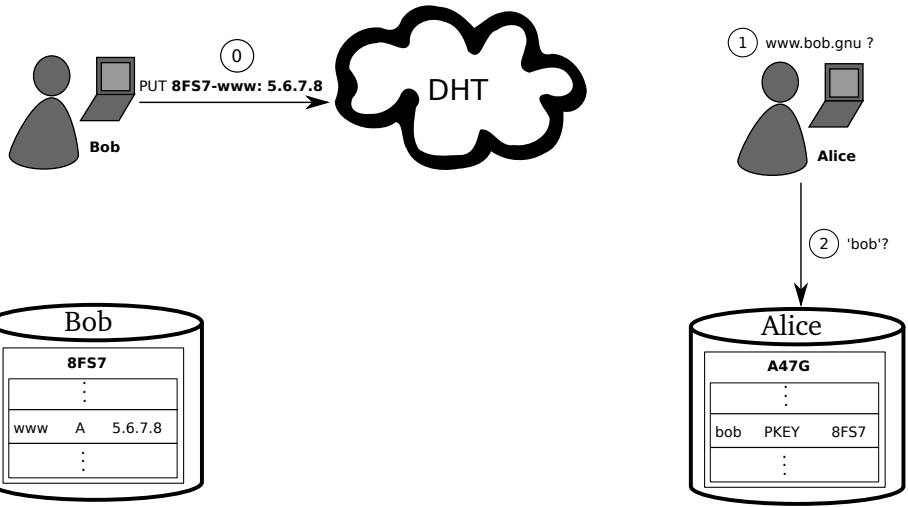

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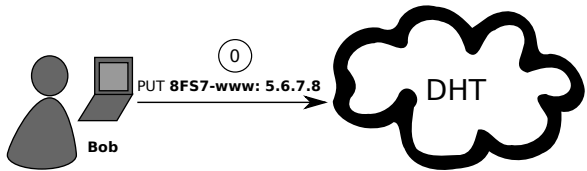

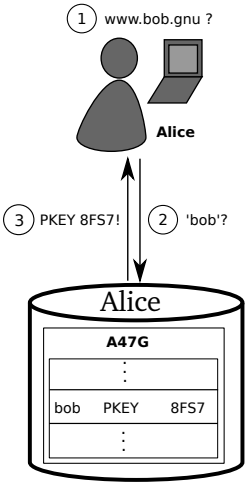

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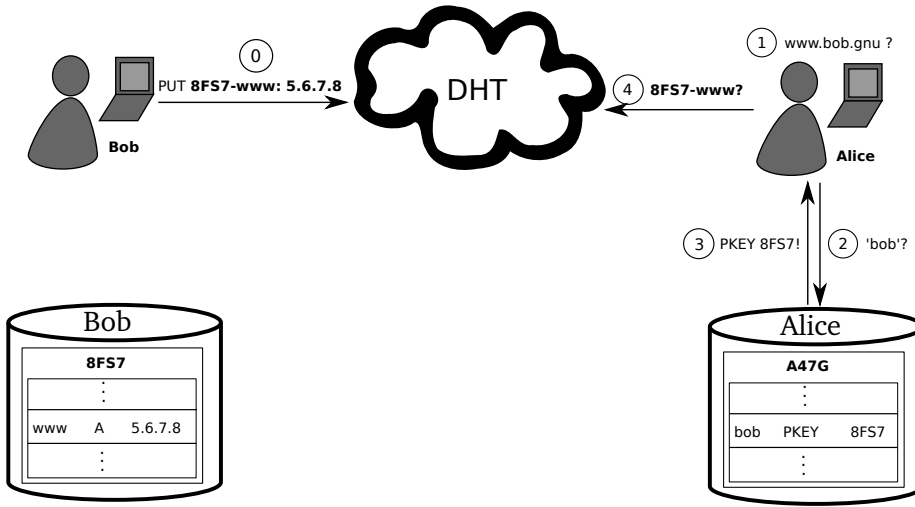

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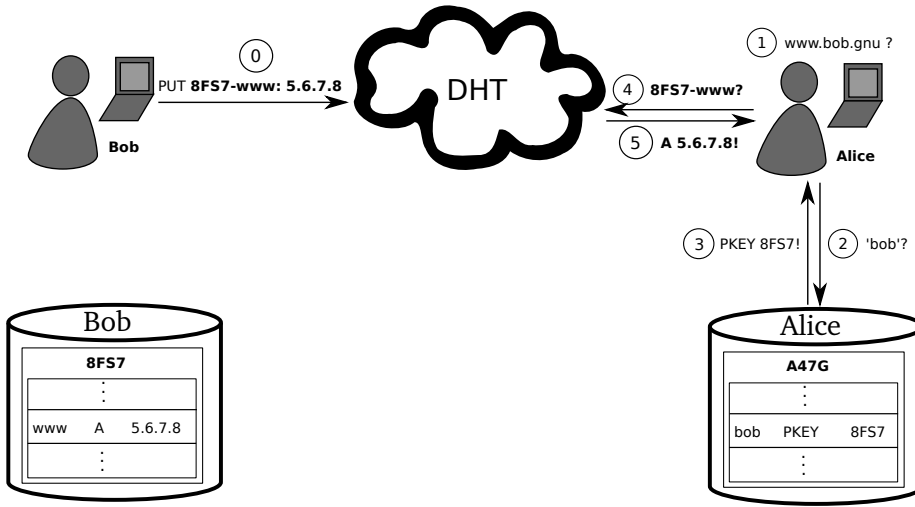
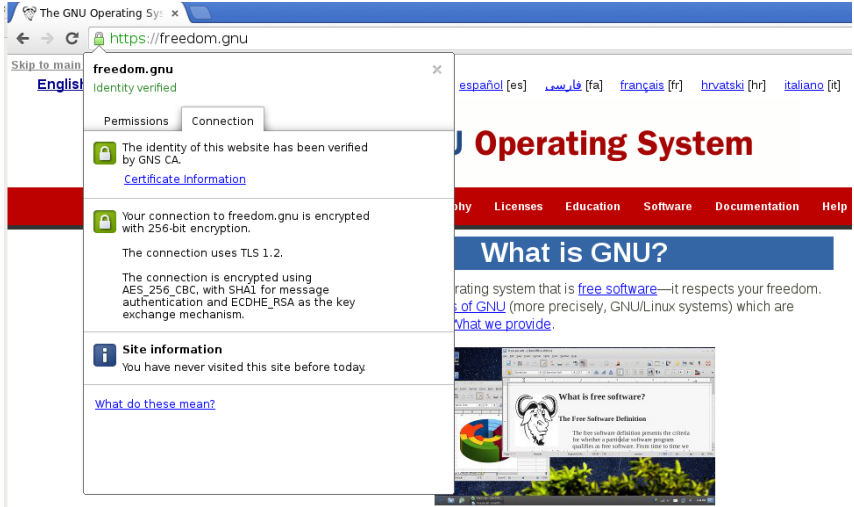
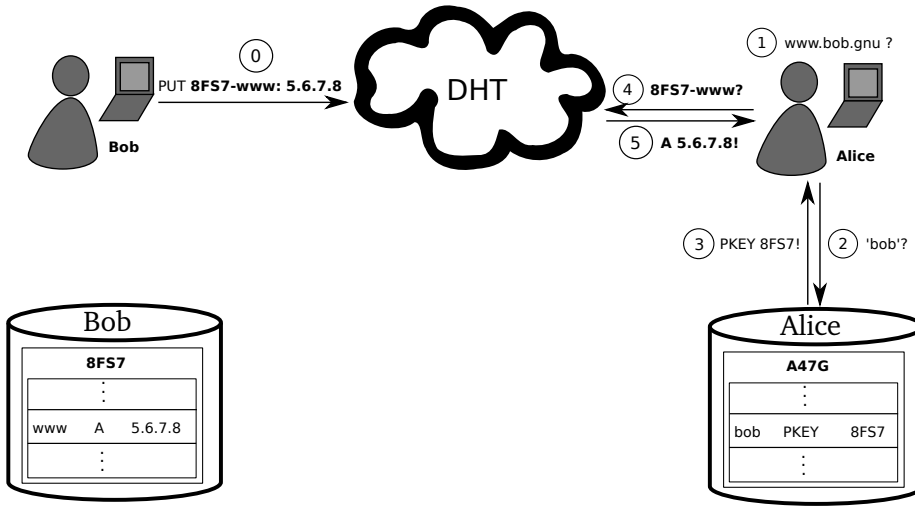


GNS as PKI (via DANE/TLSA)



Security Issue: DHT



Query Privacy: Terminology

G	generator in ECC curve, a point
n	size of ECC group, $n := \|G\|$, n prime
x	private ECC key of zone ($\in \mathbb{Z}_n$)
P	public key of zone, a point $P := xG$
l	label for record in a zone ($\in \mathbb{Z}_n$)
$R_{P,l}$	set of records for label l in zone P
$q_{P,l}$	query hash (hash code for DHT lookup)
$B_{P,l}$	block with information for label l in zone P published in the DHT under $q_{P,l}$

Query Privacy: Cryptography

Publishing B under $q_{P,I} := H(dG)$

$$h := H(I, P) \tag{1}$$
$$d := h \cdot x \mod n \tag{2}$$
$$B_{P,I} := S_d(E_{HKDF(I,P)}(R_{P,I})), dG \tag{3}$$

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Searching for I in zone P

$$h = H(I, P) \tag{4}$$
$$q_{P,I} = H(dG) = H(hxG) = H(hP) \Rightarrow \texttt{obtain } B_{P,I} \tag{5}$$
$$R_{P,I} = D_{HKDF(I,P)}(B_{P,I}) \tag{6}$$

Revocation

Revocation Basics

- Revocation certificate (RC): message signed with private key
- Peer receives new valid RC, floods to all neighbours
- All peers store all valid RCs forever

$\Rightarrow$ Expensive operation $\Rightarrow$ proof-of-work

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Revocation Magic

- Peers maybe offline during initial flood
- Network might be temporarily partitioned

$\Rightarrow$ Need to reconsile revocation sets on connect

Whenever two peers establish a P2P connection, they must compute the set union of their RC sets!

Efficient Set Union

(based on “What’s the difference? Efficient Set Reconciliation without Prior Context”, Eppstein et al., SIGCOMM’11)

- Alice and Bob have sets A and B
- The sets are very large
- . . . but their symmetric difference $\delta = |(A - B) \cup (B - A)|$ is small
- Now Alice wants to know $B - A$ (the elements she’s missing)
- . . . and Bob $A - B$ (the elements he’s missing)
- How can Alice and Bob do this efficiently?
  - w.r.t. communication and computation

Bad Solution

- Naive approach: Alice sends A to Bob, Bob sends $B - A$ back to Alice
- . . . and vice versa.

- Communication cost: $O(|A| + |B|)$:(
- Ideally, we want to do it in $O(\delta)$.
- First improvement: Don't send elements of A and B, but send/request hashes. Still does not improve complexity :(

- We need some more fancy data structure!

Bloom Filters

Constant size data structure that “summarizes” a set.

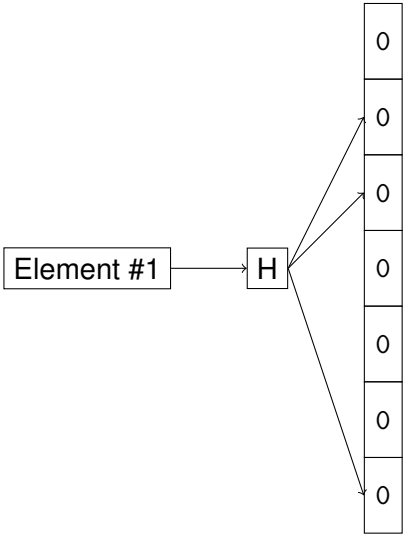
Operations:

$d = \mathit{NewBF}(\mathit{size})$ Create a new, empty bloom filter.

$\mathit{Insert}(d, e)$ Insert element e into the BF d.

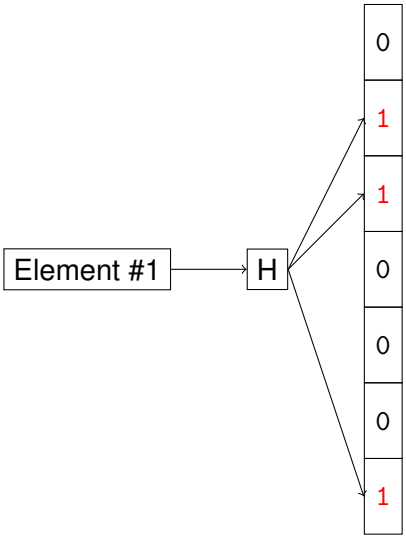
$b = \mathit{Contains}(d, e)$ Check if BF d contains element e.
$b \in \{\text{“Definitely not in set”}, \text{“Probably in set”}\}$

BF: Insert



$H(\text{Element \#1}) = (2, 3, 7)$

BF: Insert



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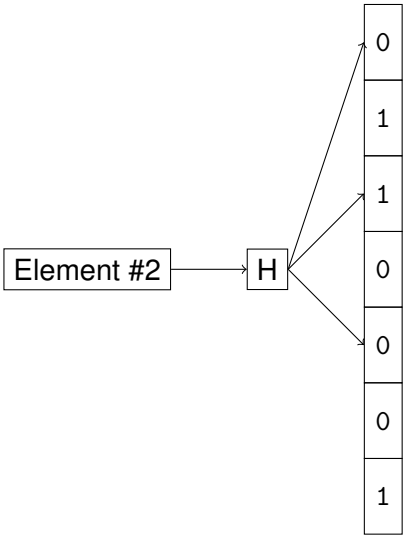
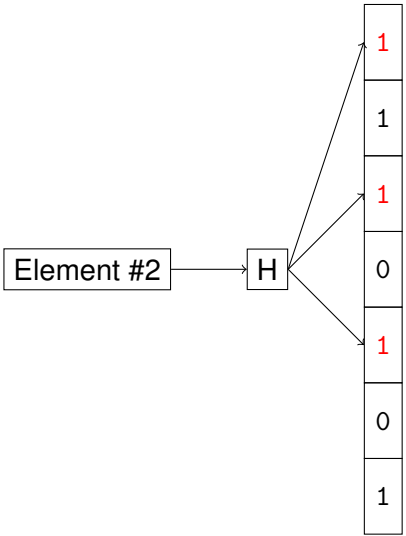


$H(\text{Element \#1}) = (2, 3, 7)$
$H(\text{Element \#2}) = (1, 3, 5)$

BF: Insert



$H(\text{Element \#1}) = (2, 3, 7)$
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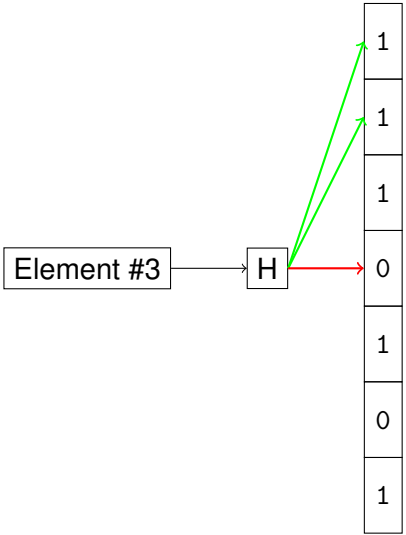
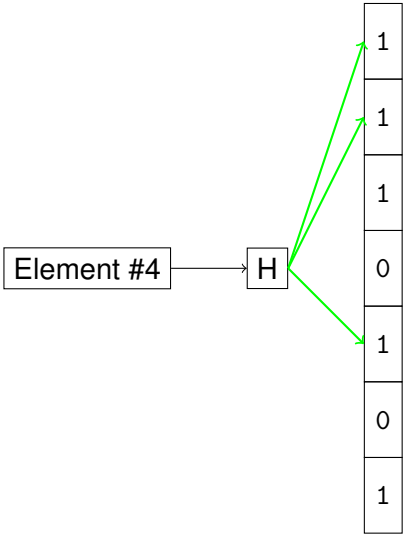
BF: Membership Test



$H(\text{Element \#1}) = (2, 3, 7)$
$H(\text{Element \#2}) = (1, 3, 5)$

BF: Membership Test (false positive)



$H(\text{Element \#1}) = (2, 3, 7)$
$H(\text{Element \#2}) = (1, 3, 5)$

Counting Bloom Filters

BF where buckets hold a **positive integer**.

Additional Operation:

Remove(d, e) Remove element from the CBF d.

$\Rightarrow$ False negatives when removing a non-existing element.

Invertible Bloom Filters

Similar to CBF, but

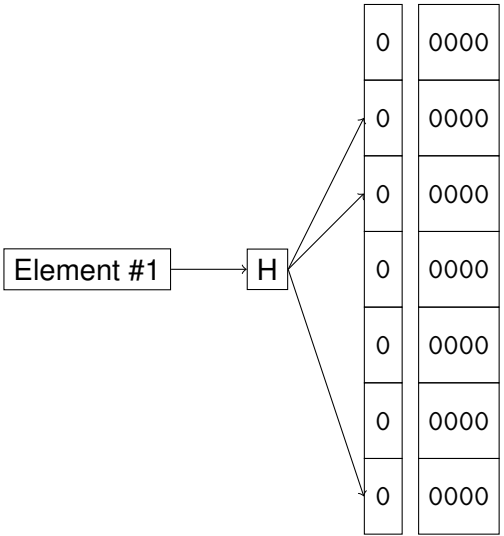
- Allow **negative counts**
- Additionaly store **(XOR-)sum of hashes** in buckets.

Additional Operations:

$(e, r) = \mathit{Extract}(d)$ Extract an element (e) from the IBF d, with result code $r \in \{\mathit{left}, \mathit{right}, \mathit{done}, \mathit{fail}\}$

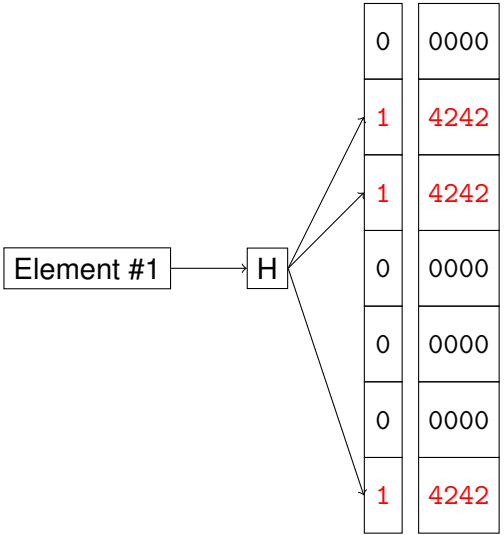
$d' = \mathit{SymDiff}(d_1, d_2)$ Create an IBF that represents the symmetric difference of d_1 and d_2.

IBF: Insert



$H(\text{Element \#1}) = (2, 3, 7)$
$H'(\text{Element \#1}) = 4242$

IBF: Insert



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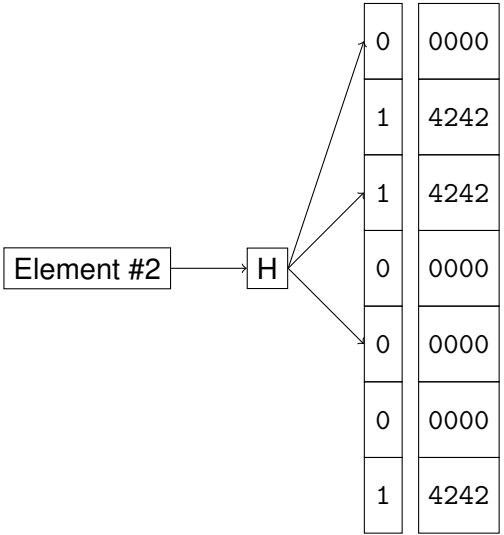
IBF: Insert



$H(\text{Element \#1}) = (2, 3, 7)$
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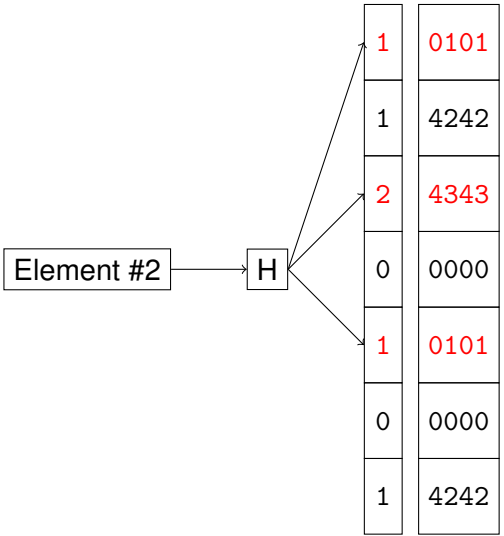
IBF: Insert



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$H'(\text{Element \#1}) = 4242$
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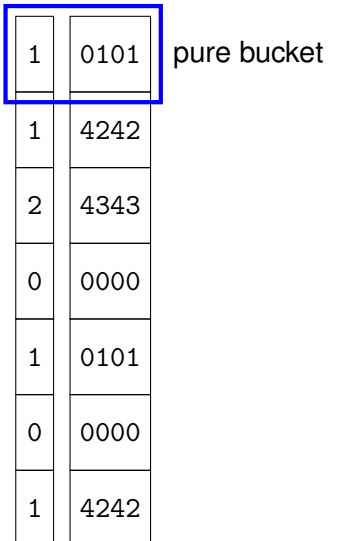
IBF: Extract



- Pure bucket $\Rightarrow$ extractable element hash
- Extraction $\Rightarrow$ more pure buckets (hopefully/probably)
- Less elements $\Rightarrow$ more chance for pure buckets

Symmetric Difference

We can directly compute the symmetric difference without extraction.

- Subtract counts
- XOR hashes

The Set Union Protocol

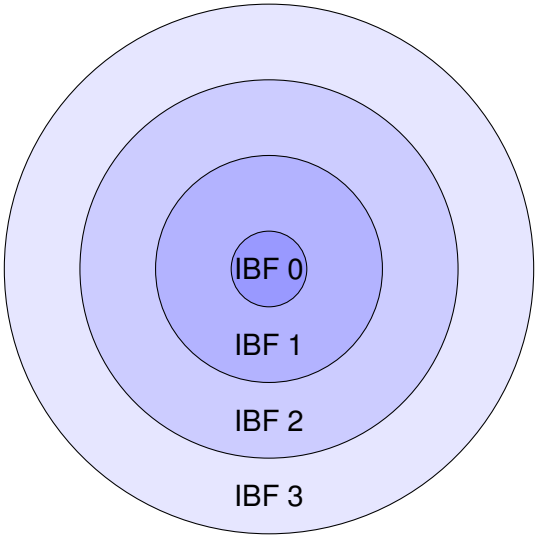
1. Create IBFs
2. Compute SymDiff
3. Extract element hashes

- Amount of communication and computation only depends on δ, not $|A| + |B|$:)
- How do we choose the initial size of the IBF?
- $\Rightarrow$ Do difference estimation first!

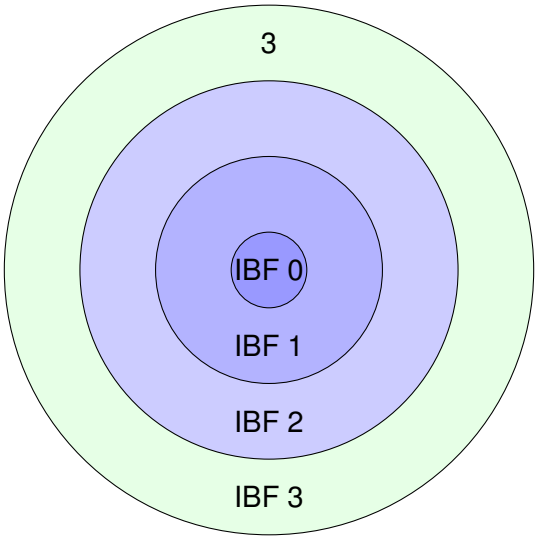
Difference Estimation

- Needed: Estimator accurate for *small* distances
- Turns out we can re-use IBFs for difference estimation
- *Sample* the set by looking at hashes, create multiple IBFs

Strata Estimator



Strata Estimator



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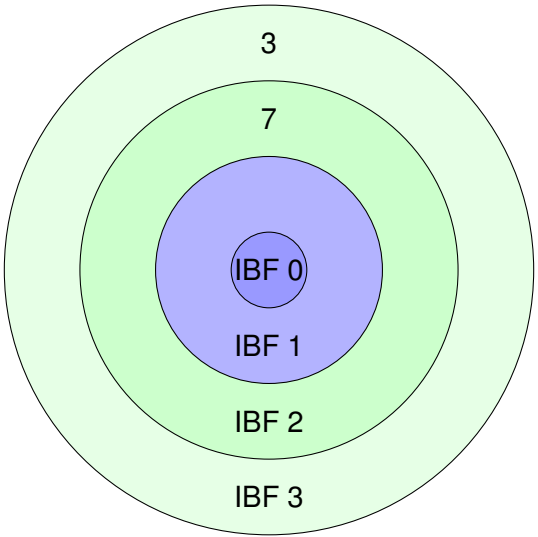

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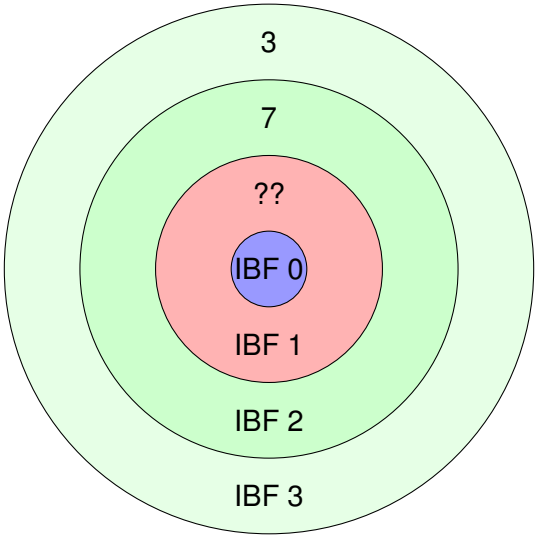
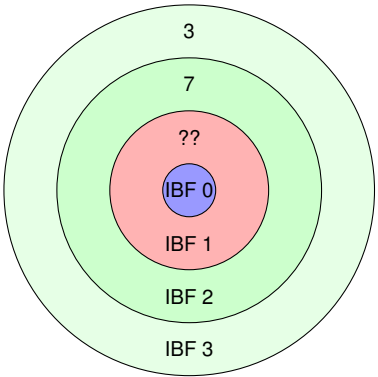


Estimation



Estimate as $(3 + 7) \cdot 2^2$.
(Number of extracted hashes scaled by expected number of elements in the remaining IBFs)

The ".zkey" pTLD

- "LABELS.*PKEY*.zkey" format
- PKEY is the public key of the zone
- Works a bit like ".onion"

$\Rightarrow$ Globally unique identifiers!

NICKnames

- “alice.bob.carol.dave.gnu” is a bit long for Eve (“.gnu”)
- Also, we need to trust Bob, Carol and Dave (for each lookup)
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- GNS adds a “NICK” record to each record set automatically
- Eve learns the “NICK”, and GNS creates “krista.short.gnu”
- Memorable, short trust path in the future! TOFU!
- Krista better pick a reasonably unique NICK.

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- Expiration time controls validity, like in DNS
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- Expiration time controls validity, like in DNS
- DHT propagation has higher delays, compared to DNS
- SHADOW is a flag in a record
- Shadow records are only valid if no other, non-expired record of the same type exists

Practical Concerns

- Name registration
- Support for browsing
- New record types
- Integration with applications
- State of the implementation

Registering a name in GNS

- Bob gives his PKEY to his **friends** via QR code
- or registers it at the **GNUnet fcfs** authority *pin.gnu* as ”bob”
- $\rightarrow$ Bob’s friends can resolve his records via *.*petname*.gnu
- $\rightarrow$ or *.bob.pin.gnu

From DNS to GNS

Names are not globally unique, but ...

... we need support for Virtual Hosting!

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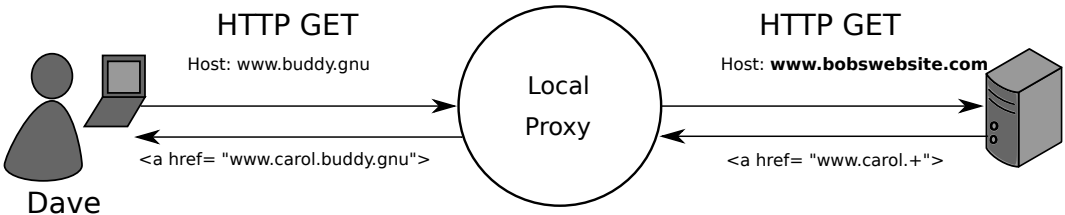
Solution: Client Side SOCKS Proxy

Legacy Hostname (LEHO) Records

LEHO records give a hint about the DNS name the server expects.



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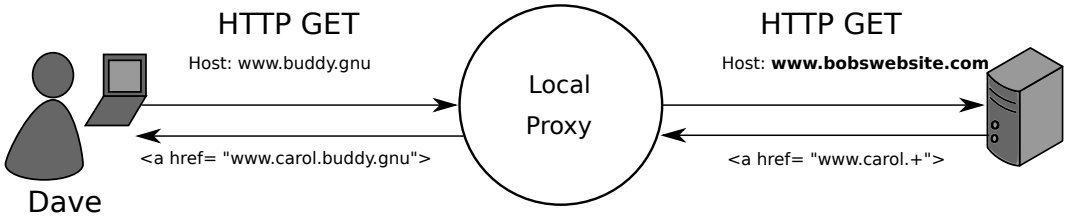

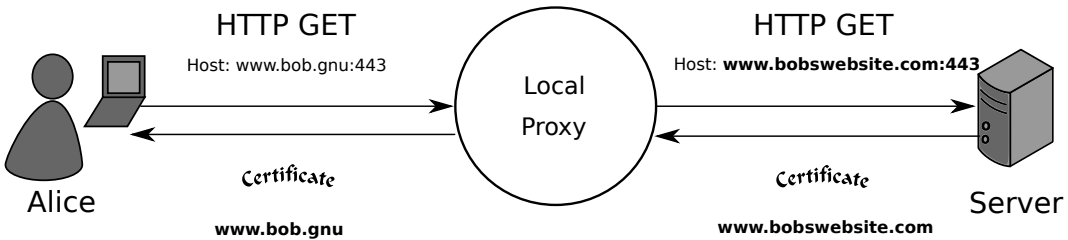

Long-Term Vision

- Integration with browser and HTTP server
- HTTP server receives “GNS-Zone: PKEY” instead of “Hostname”
- HTTP client uses “TLSA” record of GNS, instead of “LEHO”

Relative Names

- GNS records can contain ".+"
- CNAME: "server1.+"
- MX: "mail.+"
- ".+" stands for "relative to current zone"

Supporting this for links in browsers would be nice, too.

New Record Types

- PKEY: delegate to another GNS zone
- NICK: preferred names for shortening
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New Record Types

- PKEY: delegate to another GNS zone
- NICK: preferred names for shortening
- LEHO: legacy hostname
- GNS2DNS: delegate to DNS
- VPN: peers hosting TCP/IP services
- PHONE: call users using `gnunet-conversation`

DNS Delegation

- Delegate to DNS using GNS2DNS records
- GNS2DNS record specifies:
  - Name of DNS resolver (i.e. “ns1.example.com” or “piratedns.+”)
  - DNS domain to continue resolution in (i.e. “example.com” or “piratebay.org”)
- GNS will first resolve DNS resolver name to A/AAAA record
- GNS will then resolve “*left.of.gns2dns*.example.com” using DNS

VPN Delegation

- Delegates to GNUnet VPN
- VPN record specifies:
  - Identity of hosting peer (no anonymity!)
  - Service identifier (hash code)
- GNS can map VPN record to A/AAAA record of `gnunet-vpn` tunnel

PHONE service

- PHONE record specifies:
  - Identity of hosting peer (no anonymity!)
  - Line number (to support multiple phones per peer)
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ybti assembly plug

- Bart Polot will present more about Conversation
- Florian Dold will present more fun GNUnet crypto
- Julian Kirsch will present Knock

Application Integration

- SOCKS proxy (`gnunet-gns-proxy`)
- NSS plugin
- DNS packet interception (`gnunet-dns-service`)
- GNS (C) API
- GNS (IPC) protocol
- GNS command-line tool

Application Integration

```
FILE *p;
char *cmd;
char line[128];
struct in_addr ip;

if (-1 == asprintf(&cmd, "%s %s\n", "gnunet-gns -r -u", name))
  return -1;

p = popen(cmd,"r");

if (p != NULL )
{
  if (fgets( line, sizeof(line), p ) != NULL)
  {
    if (line[strlen(line)-1] == '\n')
    {
      line[strlen(line)-1] = '\0';
      if (inet_pton(af, line, &ip)))
      {
         //Do something
      }
      else
      {
        fclose (p);
        free (cmd);
        return -1;
      }
    }
  }
  ...
```

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- GNS part of GNUnet since 0.9.3
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GNS Key Exchange Party Plug

- Matthias Wachs will describe process at Lightning Talks 2
- Install GNUnet today & create private key
- Use `gnunet-bcd` to create business cards

$\Rightarrow$ Print business cards at Wau Holland tomorrow!

End

Thank you!

grothoff@in.tum.de

Get the code:

https://gnunet.org/gns